RESEARCH ARTICLE

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Subspace methods for three-parameter eigenvalue problems

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Summary

We propose subspace methods for three-parameter eigenvalue problems. Such problems arise when separation of variables is applied to separable boundary value problems; a particular example is the Helmholtz equation in ellipsoidal and paraboloidal coordinates. While several subspace methods for two-parameter eigenvalue problems exist, their extensions to a three-parameter setting seem challenging. An inherent difficulty is that, while for two-parameter eigenvalue problems, we can exploit a relation to Sylvester equations to obtain a fast Arnoldi-type method, such a relation does not seem to exist when there are three or more parameters. Instead, we introduce a subspace iteration method with projections onto generalized Krylov subspaces that are constructed from scratch at every iteration using certain Ritz vectors as the initial vectors. Another possibility is a Jacobi–Davidson-type method for three or more parameters, which we generalize from its two-parameter counterpart. For both approaches, we introduce a selection criterion for deflation that is based on the angles between left and right eigenvectors. The Jacobi–Davidson approach is devised to locate eigenvalues close to a prescribed target; yet, it often also performs well when eigenvalues are sought based on the proximity of one of the components to a prescribed target. The subspace iteration method is devised specifically for the latter task. The proposed approaches are suitable especially for problems where the computation of several eigenvalues is required with high accuracy. MATLAB implementations of both methods have been made available in the package `MultiParEig` (see http://www.mathworks.com/matlabcentral/fileexchange/47844-multipareig).

KEYWORDS

Arnoldi method, Baer wave equation, ellipsoidal wave equation, Jacobi–Davidson method, multiparameter eigenvalue problem, tensor

1 | INTRODUCTION

We consider an algebraic multiparameter eigenvalue problem of the form

$$A_{10}x_1 = \lambda_1 A_{11}x_1 + \cdots + \lambda_k A_{1k}x_1,$$
$$\vdots$$
$$A_{k0}x_k = \lambda_1 A_{k1}x_k + \cdots + \lambda_k A_{kk}x_k,$$

(1)

where $A_{ij} \in \mathbb{C}^{n_i \times n_i}$ are given matrices for $i = 1, \ldots, k$ and $j = 0, \ldots, k$. We are looking for nonzero vectors $x_i \in \mathbb{C}^{n_i}$ and a k-tuple $(\lambda_1, \ldots, \lambda_k)$ that satisfy (1). Such a k-tuple $(\lambda_1, \ldots, \lambda_k)$ is called an eigenvalue, and the tensor product $x_1 \otimes \cdots \otimes x_k$ is called the corresponding eigenvector. For more details on multiparameter eigenvalue problems, we refer to the work of Atkinson.[1]

One possible source for such problems is the separation of variables; when applied to certain separable boundary value problems (see, e.g., the works of Moon and Spencer[2] and Willatzen and Lew Yan Voon[3]), we obtain a system of k linear ordinary differential equations of the form

$$p_j(x_j) \, y_j''(x_j) + q_j(x_j) \, y_j'(x_j) + r_j(x_j) \, y_j(x_j) = \sum_{\ell=1}^{k} \lambda_\ell \, s_{j\ell}(x_j) \, y_j(x_j), \quad j = 1, \ldots, k, \tag{2}$$

where $x_j \in [a_j, b_j]$, together with appropriate boundary conditions. We are interested in a k-tuple $(\lambda_1, \ldots, \lambda_k)$ and nontrivial functions $y_1, \ldots, y_k$ such that Equation (2) and the boundary conditions are satisfied. For more details on systems of the form (2), we refer to the work of Atkinson and Mingarelli[4]; see also Section 2.

By discretizing (2), we obtain a problem of the form (1). This approach is used in the work of Plestenjak et al.[5] to find numerical solutions for several separable boundary value problems and improve previous results from the literature. Specifically, spectral collocation is used in the aforementioned work[5] for the discretization, which gives rise to relatively small matrices and accurate results. While several suitable numerical methods for the case $k = 2$ exist (see, e.g., the work of Plestenjak et al.[5] and the references therein), available feasible numerical methods for $k \geq 3$ are limited to problems with very small matrices, which means that, even by using spectral collocation, we cannot obtain many accurate eigenvalues of (2). We introduce new variants of numerical methods for three-parameter eigenvalue problems that exceed the above limitations and can be applied to problems with larger matrices. This allows us to solve efficiently and accurately several three-parameter eigenvalue problems of the form (2), which we demonstrate in numerical examples.

Let S_k denote the set of permutations of the set $\{1, \ldots, k\}$, and let $\mathrm{sgn}(\sigma)$ be the sign of a permutation $\sigma \in S_k$. By introducing the $k \times k$ operator determinants

$$\Delta_0 := \begin{vmatrix} A_{11} & \cdots & A_{1k} \\ \vdots & & \vdots \\ A_{k1} & \cdots & A_{kk} \end{vmatrix}_{\otimes} = \sum_{\sigma \in S_k} \mathrm{sgn}(\sigma) \, A_{1\sigma_1} \otimes A_{2\sigma_2} \otimes \cdots \otimes A_{k\sigma_k}, \tag{3}$$

where $\otimes$ denotes the Kronecker product, and, similarly,

$$\Delta_i := \begin{vmatrix} A_{11} & \cdots & A_{1,i-1} & A_{10} & A_{1,i+1} & \cdots & A_{1k} \\ \vdots & & \vdots & \vdots & \vdots & & \vdots \\ A_{k1} & \cdots & A_{k,i-1} & A_{k0} & A_{k,i+1} & \cdots & A_{kk} \end{vmatrix}_{\otimes} \tag{4}$$

for $i = 1, \ldots, k$, we obtain matrices $\Delta_0, \ldots, \Delta_k$ of size $(n_1 \cdots n_k) \times (n_1 \cdots n_k)$. If Δ_0 is nonsingular, then the matrices $\Delta_0^{-1}\Delta_1, \ldots, \Delta_0^{-1}\Delta_k$ commute, and (1) is equivalent to a system of generalized eigenvalue problems

$$\Delta_j z = \lambda_j \Delta_0 z, \quad j = 1, \ldots, k$$

for $z = x_1 \otimes \cdots \otimes x_k$ (for details, see, e.g., the work of Atkinson[1]). This relation enables one to use standard numerical methods for generalized eigenvalue problems if the Δ-matrices are not too large. However, when spectral methods are used to discretize (2), then, in practice, even for $k = 2$, the Δ-matrices might be so large that it is not efficient, or even not feasible, to compute all of the eigenvalues. Fortunately, for various applications, the retrieval of several eigenvalues closest to a prescribed target is sufficient. In some other cases, eigenvalues $(\lambda_1, \ldots, \lambda_k)$ such that a prescribed component among $\lambda_1, \ldots, \lambda_k$ is close to a given target σ are of interest. For instance, when we apply separation of variables to the k-dimensional Helmholtz equation $\nabla^2 u + \omega^2 u = 0$, usually, only one of the parameters $\lambda_1, \ldots, \lambda_k$ is related to the eigenfrequency ω (see Section 2 for more details). If we assume, without loss of generality, that λ_k is relevant to the problem and we are interested in the first low-frequency modes for the Helmholtz equation, then we are looking for eigenvalues with the smallest value of $|\lambda_k|$.

1.1 | Overview

Jacobi–Davidson-type methods have been proposed for the two-parameter eigenvalue problem in the works of Hochstenbach et al.[6] and Hochstenbach and Plestenjak[7] to compute a few eigenvalues closest to a prescribed target. When eigenvalues $(\lambda_1, \ldots, \lambda_k)$ with the smallest $|\lambda_k|$ are sought, subspace iteration or an Arnoldi iteration operating directly on

$\Delta_k z = \lambda_k \Delta_0 z$ appears more appropriate. Such ideas have been explored well in the two-parameter eigenvalue setting and applied for the solution of various separable boundary value problems.[5,8] This success is mostly due to the fact that linear systems of the form $\Delta_2 w = \Delta_0 v$ for a given v can be expressed as Sylvester equations involving the matricizations of the vectors v and w and, thus, can be solved efficiently at a cost of $O(n_1^3 + n_2^3)$. An underlying difficulty is that such a Sylvester equation representation is not known for the linear system $\Delta_k w = \Delta_0 v$ when $k \geq 3$.

The main contributions of this work are a Jacobi–Davidson method in Section 4.2 and an inexact subspace iteration method with Ritz projections in Section 4.5 for three-parameter eigenvalue problems. The Jacobi–Davidson method is inspired by earlier works,[6,7] but new ingredients are also put in use. For instance, a Newton method–based tensor Rayleigh quotient iteration (TRQI) is incorporated to speed up convergence. Numerical experiments indicate that the proposed Jacobi–Davidson method is effective in extracting both the eigenvalues closest to a prescribed target and the eigenvalues $(\lambda_1, \lambda_2, \lambda_3)$ whose λ_3 components are closest to a prescribed target. On the other hand, inexact subspace iteration, which operates directly on the generalized eigenvalue problem $\Delta_3 z = \lambda_3 \Delta_0 z$, is tailored to compute eigenvalues with their λ_3 components closest to a prescribed target. Instead of solving a linear system of the form $\Delta_3 w = \Delta_0 v$ for the unknown w, it projects the full problem onto certain generalized Krylov subspaces that are restarted at every iteration with selected Ritz vectors. We especially aim at problems where the computation of several extreme eigenvalues is required with high accuracy. Both of the proposed Jacobi–Davidson method and inexact subspace iteration are well suited to deal with such problems.

1.2 | Outline

We start with two particular applications giving rise to three-parameter eigenvalue problems in Section 2; this is followed by a brief review of subspace iteration approaches for the two-parameter case in Section 3. In particular, efficient solutions of the linear system $\Delta_2 w = \Delta_0 v$ with or without projections via their Sylvester equation characterization facilitate these approaches.

The main body is Section 4, which introduces iterative methods for the extraction of a few targeted eigenvalues of a three-parameter eigenvalue problem. A Jacobi–Davidson method is proposed in Section 4.2. The difficulty intrinsic to applying a Krylov subspace method directly to $\Delta_3 z = \lambda_3 \Delta_0 z$ is pointed out in Section 4.3. Consequently, in Section 4.4, a subspace iteration method that does not work on the full linear systems, but rather solves their projections onto Krylov subspaces, is described. The downside of this approach is that, in every iteration, it requires low-rank third-order tensor approximations for the solutions of the linear systems. Finally, an efficient Krylov subspace–based subspace iteration is proposed in Section 4.5, which employs the projection ideas in Section 4.4, but removes the need for low-rank tensor approximations.

Section 5 is devoted to extensive numerical experiments. In particular, we illustrate how the proposed Jacobi–Davidson and subspace iteration methods perform on the three-parameter eigenvalue problems resulting from the applications in Section 2, as well as on a random synthetic example.

2 | MOTIVATION

We give two applications that lead to three-parameter eigenvalue problems of the form (2). They concern the separation of variables applied to the Helmholtz equation

$$\nabla^2 u + \omega^2 u = 0 \tag{5}$$

in ellipsoidal and paraboloidal coordinates.

2.1 | Ellipsoidal wave equations

If we aim to compute eigenfrequencies of an ellipsoidal body with a fixed boundary, then we have to solve the Helmholtz equation (5) over the ellipsoid

$$\Omega := \left\{ (x, y, z) \in \mathbb{R}^3 \mid (x/x_0)^2 + (y/y_0)^2 + (z/z_0)^2 \leq 1 \right\}$$

subject to the Dirichlet boundary condition $u|_{\partial\Omega} = 0$. Here, x_0, y_0, z_0 correspond to the radii of the semi-axes of the ellipsoid and satisfy $z_0 > y_0 > x_0 > 0$. A numerical approach has been proposed in the work of Willatzen and Lew Yan Voon[9]

(see also the works of Plestenjak et al.[5] and Levitina[10]); here, we give an outline of how it leads to a three-parameter eigenvalue problem.

The Helmholtz equation is separable in ellipsoidal coordinates (ξ_1, ξ_2, ξ_3),[2] a natural choice for the region Ω. Formally, there exist functions $X_1(\xi_1), X_2(\xi_2)$, and $X_3(\xi_3)$ such that the solution can be written as

$$u\left(x(\xi_1, \xi_2, \xi_3),\ y(\xi_1, \xi_2, \xi_3),\ z(\xi_1, \xi_2, \xi_3)\right) = X_1(\xi_1)X_2(\xi_2)X_3(\xi_3).$$

Exploiting the separability property above and expressing the Helmholtz equation in ellipsoidal coordinates, we obtain three ordinary differential equations

$$t_j(t_j - 1)(t_j - c)\tilde{X}_j'' + \frac{1}{2}\left(3t_j^2 - 2(1 + c)t_j + c\right)\tilde{X}_j' + \left(\lambda + \mu t_j + \eta t_j^2\right)\tilde{X}_j = 0, \quad j = 1, 2, 3,$$

where $c = a^2/b^2$, $a = (z_0^2 - x_0^2)^2$, $b = (z_0^2 - y_0^2)^2$, $t_j = \xi_j^2/b^2$, $\tilde{X}_j(t_j) := X_j(\xi_j(t_j))$, and the elliptical coordinates satisfy $z_0 > \xi_1 > a > \xi_2 > b > \xi_3 > 0$. The three differential equations are coupled by the scalars λ, μ, and η, but only $\eta = \omega^2 b^2/4$ is related to the eigenfrequency ω. The function $\tilde{X}_j(t_j)$ above is of the form

$$\tilde{X}_j(t_j) = t_j^{\rho/2}(t_j - 1)^{\sigma/2}(t_j - c)^{\tau/2}F_j(t_j),$$

where $F_j(t_j)$ is an integral function of t_j, and ρ, σ, τ can take values 0 or 1. For each one of the eight possible configurations for (ρ, σ, τ), we deduce the system of ordinary differential equations

$$t_j(t_j - 1)(t_j - c)F_j'' + \frac{1}{2}\left(k_2 t_j^2 - 2k_1 t_j + k_0\right)F_j' + \left(\lambda - \lambda_0 + (\mu + \mu_0)t_j + \eta t_j^2\right)F_j = 0, \quad j = 1, 2, 3, \tag{6}$$

with

$$\lambda_0 = \frac{1}{4}\left[(\rho + \tau)^2 + (\rho + \sigma)^2 c\right], \quad \mu_0 = \frac{1}{4}(\rho + \sigma + \tau)(\rho + \sigma + \tau + 1),$$

$$k_0 = (2\rho + 1)c, \quad k_1 = (1 + \rho)(1 + c) + \tau + \sigma c, \quad k_2 = 2(\rho + \sigma + \tau) + 3.$$

The boundedness conditions at singular points and the Dirichlet condition on the boundary of the ellipsoid give rise to the following boundary conditions:

$$F_1\left(z_0^2/b^2\right) = 0,$$

$$(k_2 c^2 - k_1 c + k_0)F_j'(c) + 2(\lambda - \lambda_0 + (\mu + \mu_0)c + \eta c^2)F_j(c) = 0 \quad \text{for} \quad j = 1, 2,$$

$$(k_2 - 2k_1 + k_0)F_j'(1) + 2(\lambda - \lambda_0 + \mu + \mu_0 + \eta)F_j(1) = 0 \quad \text{for} \quad j = 2, 3,$$

$$k_0 F_3'(0) + 2(\lambda - \lambda_0)F_3(0) = 0.$$

This example will be solved numerically in Section 5.1 more accurately than in the work of Plestenjak et al.[5] as the new methods can deal with larger matrices coming from finer discretizations.

2.2 | Baer wave equations

The Helmholtz equation (5) is also separable in paraboloidal coordinates (ξ_1, ξ_2, ξ_3), which are related to the Cartesian coordinates by (see, e.g., the works of Moon and Spencer[2] and Duggen et al.[11])

$$x^2 = 4(c - b)^{-1}(b - \xi_1)(b - \xi_2)(b - \xi_3),$$

$$y^2 = 4(b - c)^{-1}(c - \xi_1)(c - \xi_2)(c - \xi_3),$$

$$z = \xi_1 + \xi_2 + \xi_3 - b - c,$$

where $-\infty < \xi_1 < c < \xi_2 < b < \xi_3 < \infty$ and $c < b$ are the parameters of the paraboloidal coordinate system. A constant surface $\xi_1 = \gamma$, where $\gamma < c$, represents an upward opening elliptic paraboloid that intersects the z-axis at $z = \gamma$, whereas a constant surface $\xi_3 = \beta$, where $b < \beta$, represents a downward opening elliptic paraboloid that intersects the z-axis at $z = \beta$.

In Section 5.2, we will consider the solution of the Helmholtz equation with a fixed boundary on a domain bounded by the two elliptic paraboloids $\gamma = 0$ and $\beta = 5$, as well as for the choices of $c = 1$ and $b = 3$ (see Figure 1).
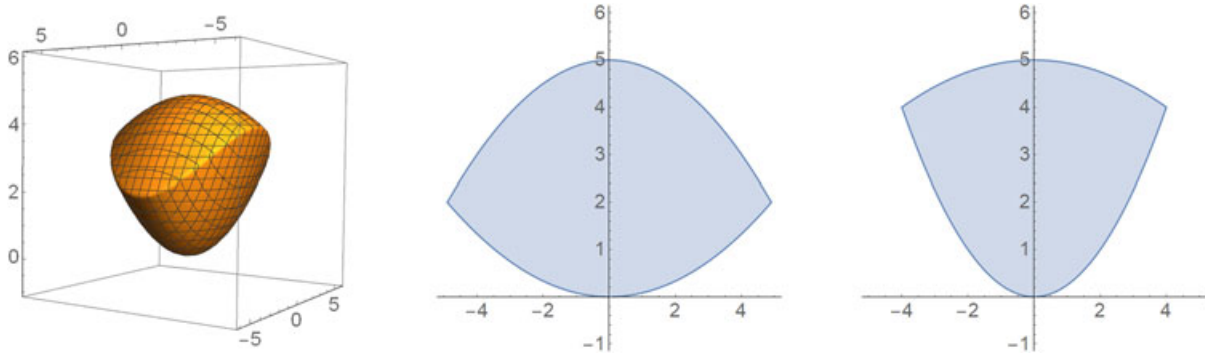
FIGURE 1 Region bounded in paraboloidal coordinates by elliptical paraboloids $\xi_1 = 0$ and $\xi_3 = 5$ (left) and its intersection with the xz-plane (middle) and the yz-plane (right)

We use separation of variables. The solution of (5) has the form $u = X_1(\xi_1) X_2(\xi_2) X_3(\xi_3)$,[2] where X_1, X_2, X_3 satisfy the system of Baer wave differential equations given by

$$(\xi_j - b)(\xi_j - c)X_j'' + \frac{1}{2}\left(2\xi_j - (b+c)\right)X_j' + \left(\lambda + \mu\xi_j + \eta\xi_j^2\right)X_j = 0, \qquad j = 1, 2, 3, \tag{7}$$

and ξ_1, ξ_2, ξ_3 are such that $\gamma < \xi_1 < c < \xi_2 < b < \xi_3 < \beta$. In the equations above, $\eta = \omega^2$ is related to the eigenfrequency, whereas parameters λ and μ result from the separation. Equation (7) has regular singularities at b and c and an irregular singularity at infinity. The exponents at the finite singularities are 0 and $1/2$. Therefore, it is possible to write the solution of (7) as

$$X_i(\xi_i) = (\xi_i - b)^{\rho/2}(\xi_i - c)^{\sigma/2} F_i(\xi_i), \tag{8}$$

where $F_i(\xi_i)$ is an integral function of ξ_i, and ρ, σ can be either 0 or 1, leading to four possible configurations.

For a particular (ρ, σ) configuration, by plugging (8) into (7), we obtain the system

$$(\xi_j - b)(\xi_j - c)F_j'' + \frac{1}{2}(k_1\xi_j - k_0)F_j' + \left(\lambda - \lambda_0 + \mu\xi_j + \eta\xi_j^2\right) F_j = 0, \quad j = 1, 2, 3 \tag{9}$$

of differential equations, where

$$k_1 = 2(1 + \rho + \sigma), \quad k_0 = (1 + 2\sigma)b + (1 + 2\rho)c, \quad \lambda_0 = -\frac{1}{4}(\rho + \sigma + 2\rho\sigma).$$

The boundedness conditions at singular points and the Dirichlet condition on the boundary of the domain yield the following boundary conditions:

$$F_1(\gamma) = 0,$$
$$\frac{1}{2}(k_1 c - k_0)F_j'(c) + (\lambda - \lambda_0 + \mu c + \eta c^2)F_j(c) = 0 \quad \text{for} \quad j = 1, 2,$$
$$\frac{1}{2}(k_1 b - k_0)F_j'(b) + (\lambda - \lambda_0 + \mu b + \eta b^2)F_j(b) = 0 \quad \text{for} \quad j = 2, 3, \tag{10}$$
$$F_3(\beta) = 0.$$

We will present some numerical experiments with these examples in Section 5.

3 | TWO PARAMETERS

In this section, we consider (1) for the case $k = 2$, but, to ease the notation, set $A_j := A_{j0}$, $B_j := A_{j1}$, $C_j := A_{j2}$ for $j = 1, 2$, as well as $\lambda := \lambda_1$, $\mu := \lambda_2$. A quick overview of the ideas in the work of Meerbergen and Plestenjak[8] is presented next. As we shall see in the subsequent section, most of these ideas for the two-parameter case cannot be generalized to more than two parameters.

Recall that if $\Delta_0 = B_1 \otimes C_2 - C_1 \otimes B_2$ is nonsingular, then the two-parameter problem at hand is equivalent to a coupled pair of generalized eigenvalue problems $\Delta_1 z = \lambda \Delta_0 z$ and $\Delta_2 z = \mu \Delta_0 z$ for $z = x_1 \otimes x_2$, where $\Delta_1 = A_1 \otimes C_2 - C_1 \otimes A_2$

and $\Delta_2 = B_1 \otimes A_2 - A_1 \otimes B_2$. Suppose that we are looking for the eigenvalues (λ, μ) with the smallest value of $|\mu|$, and let us assume that $n_1 n_2$ is so large that we cannot efficiently compute all eigenvalues of the generalized eigenvalue problem

$$\Delta_2 z = \mu \Delta_0 z. \tag{11}$$

Next, we discuss two alternative numerical approaches for this setting: Krylov subspace methods and a subspace iteration.

3.1 | Krylov subspace methods

If $n_1 n_2$ is not too large, then we can apply a Krylov subspace method to (11), for instance, the implicitly restarted Arnoldi[12] or the Krylov–Schur method.[13] As we are interested in the smallest values of $|\mu|$, we want to build an orthogonal basis for the Krylov subspace $\mathcal{K}_k(\Delta_2^{-1}\Delta_0, v_0)$, which means that in each step, we have to compute a matrix–vector product with the matrix $\Delta_2^{-1}\Delta_0$. The key observation to perform this multiplication efficiently is its connection with a Sylvester equation. That is, the expression $w = \Delta_2^{-1}\Delta_0 v$ can be rearranged as

$$(B_1 \otimes A_2 - A_1 \otimes B_2)w = (B_1 \otimes C_2 - C_1 \otimes B_2)v. \tag{12}$$

Using the vectorization operator

$$X := [X_1 \cdots X_q] \in \mathbb{C}^{p \times q}, \ X_1, \ldots, X_q \in \mathbb{C}^p \ \mapsto \ \mathrm{vec}(X) := \begin{bmatrix} X_1^T \cdots X_q^T \end{bmatrix}^T \in \mathbb{C}^{pq}$$

and the identity $(B \otimes A)\mathrm{vec}(X) = \mathrm{vec}(AXB^T)$, we can write the linear system in (12) as

$$A_2 W B_1^T - B_2 W A_1^T = C_2 V B_1^T - B_2 V C_1^T,$$

where W and V are matrices such that $\mathrm{vec}(W) = w$ and $\mathrm{vec}(V) = v$. If we assume that B_1 and B_2 are nonsingular, then this is equivalent to the Sylvester equation

$$B_2^{-1} A_2 W - W A_1^T B_1^{-T} = M, \tag{13}$$

where $M = B_2^{-1} C_2 V - V C_1^T B_1^{-T}$. As the Sylvester equation in (13) can be solved in $\mathcal{O}(n_1^3 + n_2^3)$ operations using, for example, the Bartels–Stewart method,[14] this is much more efficient than forming Δ_2 and Δ_0 explicitly, and then solving $\Delta_2 w = \Delta_0 v$, which typically requires $\mathcal{O}(n_1^3 n_2^3)$ operations.

If $n_1 n_2$ is even larger, we can neither store many vectors from $\mathcal{K}_k(\Delta_2^{-1}\Delta_0, v_0)$ fully nor perform exact computations with them efficiently. In the limit (as we keep multiplying with $\Delta_2^{-1}\Delta_0$), v and w in (12) are collinear to the dominant eigenvector of the form $z = x \otimes y$, and the corresponding matrices V and W have both rank 1. Hence, the right-hand side $M = B_2^{-1} C_2 V - V C_1^T B_1^{-T}$ of the Sylvester equation (13) is nearly of rank 2 at later iterations, whereas its solution W has almost rank 1. In this case, it is possible to benefit from an approximate low-rank solver for the Sylvester equation; see, for example, the work of Meerbergen and Plestenjak[8] that makes use of an approximate Krylov subspace solver due to the work of Hu and Reichel.[15]

The main idea of the Hu–Reichel method is as follows. Suppose that the Sylvester equation

$$AX - XB = C \tag{14}$$

is such that C has low rank, and additionally suppose that the solution X is expected to have low rank (in practice, it is enough that both C and X are close to low-rank matrices). If $C \approx FG^T$, where F and G have a few columns, then matrices Q_A and Q_B whose columns form orthonormal bases for the Krylov subspaces $\mathcal{K}_r(A, F)$ and $\mathcal{K}_r(B^T, G)$ are built. An approximate solution for (14) is then given by $X = Q_A Y Q_B^H$, where the matrix Y is the solution of the small-scale projected Sylvester equation

$$Q_A^H A Q_A Y - Y Q_B^H B Q_B = Q_A^H C Q_B.$$

For further details and various other numerical approaches for large-scale Sylvester equations, we refer to the survey paper by Simoncini[16] and the references therein.

3.2 | Subspace iteration

The subspace iteration starts with a matrix $Z_0 \in \mathbb{C}^{n_1 n_2 \times p}$, such that $Z_0^H Z_0 = I$, where $p \ll n_1 n_2$. In each step, first, p linear systems $\Delta_2 W_{k+1} = \Delta_0 Z_k$ are solved, and then, the columns of W_{k+1} are orthonormalized into Z_{k+1}. As k goes to infinity, $W_{k+1}^H Z_k$ under mild conditions converges to an upper triangular matrix with eigenvalues of the pencil (Δ_2, Δ_0) on

its diagonal. As discussed in the previous subsection, the linear systems $\Delta_2 W_{k+1} = \Delta_0 Z_k$ can be expressed as a set of p Sylvester equations, each one of which generically possesses a low-rank structure when converging. Consequently, instead of working with full vectors in the columns of the matrices Z_k and W_k, we rather use their low-rank approximations. We express the columns of Z_k as $z_i^{(k)} = \text{vec}(U_k D_i^{(k)} V_k^T)$ for $i = 1, \ldots, p$, where $U_k \in \mathbb{C}^{n_1 \times \ell}$ and $V_k \in \mathbb{C}^{n_2 \times \ell}$ have orthonormal columns and $D_i^{(k)}$ is an $\ell \times \ell$ core matrix, where $\ell \geq p$. Instead of solving $\Delta_2 w_i^{(k+1)} = \Delta_0 z_i^{(k)}$ exactly, we solve this only approximately and obtain a low-rank approximation for the matricization of $w_i^{(k+1)}$ by means of the Hu–Reichel method. Specifically, we search for $w_i^{(k+1)}$ in the space

$$\mathcal{K}_r\left(B_1^{-1}A_1, G\right) \otimes \mathcal{K}_r\left(B_2^{-1}A_2, F\right) \tag{15}$$

for a modest r, where $F = [\, B_2^{-1}C_2 U_k \;\; U_k \,]$ and $G = [\, B_1^{-1}C_1 V_k \;\; V_k \,]$ (i.e., setting $C \approx FG^T$ in the Sylvester equation (14) equal to M as in (13) yields these choices for F and G). For each $w_i^{(k+1)}$, $i = 1, \ldots, p$, we solve a small projected Sylvester equation. If the columns of $\widetilde{V}_k$ and $\widetilde{U}_k$ form orthonormal bases for $\mathcal{K}_r(B_1^{-1}A_1, G)$ and $\mathcal{K}_r(B_2^{-1}A_2, F)$, respectively, then $w_i^{(k+1)} = \text{vec}(\widetilde{U}_k Y_i^{(k)} \widetilde{V}_k^T)$ for some $Y_i^{(k)}$. For the construction of U_{k+1} and V_{k+1}, we note that all vectors $w_1^{(k+1)}, \ldots, w_p^{(k+1)}$ in the next step lie in (15). This inspired a new method in the work of Meerbergen and Plestenjak[8] called subspace iteration with Arnoldi expansion. The essential idea is to compute matrices $\widetilde{V}_k, \widetilde{U}_k$ whose columns form orthonormal bases for $\mathcal{K}_r(B_1^{-1}A_1, G), \mathcal{K}_r(B_2^{-1}A_2, F)$ and then to compute the Ritz values with the smallest values of $|\tau|$ as well as the Ritz vectors from the projected small-scale two-parameter eigenvalue problem

$$\widetilde{V}_k^H A_1 \widetilde{V}_k c = \sigma \widetilde{V}_k^H B_1 \widetilde{V}_k c + \tau \widetilde{V}_k^H C_1 \widetilde{V}_k c$$
$$\widetilde{U}_k^H A_2 \widetilde{U}k d = \sigma \widetilde{U}_k^H B_2 \widetilde{U}_k d + \tau \widetilde{U}_k^H C_2 \widetilde{U}_k d.$$

From ℓ such Ritz vectors, which are all decomposable, we form the new subspaces $\text{span}\{U_{k+1}\}$ and $\text{span}\{V_{k+1}\}$ for the next step.

4 | THREE PARAMETERS

Let us now focus on three-parameter eigenvalue problems, which are of the form (1) for $k = 3$. Similarly to the previous section, to ease the notation, we let $A_j := A_{j0}, B_j := A_{j1}, C_j := A_{j2}, D_j := A_{j3}$ for $j = 1, 2, 3$, where A_{jk} are as in (1), and $\lambda := \lambda_1, \mu := \lambda_2, \eta := \lambda_3$. In this three-parameter eigenvalue setting, we are seeking the eigenvalues with the smallest values of $|\eta|$. They correspond to the eigenvalues of the generalized eigenvalue problem

$$\Delta_3 z = \eta \Delta_0 z \tag{16}$$

with the smallest values of $|\eta|$, provided Δ_0 is nonsingular, where Δ_0 and Δ_3 denote matrices involving third-order tensors defined by (3) and (4). If such an eigenvalue is simple, then the corresponding eigenvector z is decomposable and can be expressed as $z = x_1 \otimes x_2 \otimes x_3$.

4.1 | Using full Δ-matrices

The first option is to explicitly form the matrices Δ_0 and Δ_3 and then use the QZ algorithm (or any other numerical method) to compute the eigenvalues of (16). As the size of the matrices Δ_0 and Δ_3 is $n_1 n_2 n_3 \times n_1 n_2 n_3$, this is efficient only when $n_1 n_2 n_3$ is small. This approach becomes prohibitively expensive even for modest values of n_1, n_2, and n_3.

4.2 | Jacobi–Davidson-type method

Methods of Jacobi–Davidson type have been developed for two-parameter eigenvalue problems in the works of Hochstenbach et al.[6] and Hochstenbach and Plestenjak.[7,17] As long as we are able to solve a small projected problem efficiently, the method can be generalized to multiparameter eigenvalue problems with three or more parameters.

Inspired by its two-parameter counterpart in the work of Hochstenbach and Plestenjak,[7] we give a brief description of a Jacobi–Davidson-type method for a three-parameter eigenvalue problem in Algorithm 1. In the description, $\|r_j\|$ represents the 2-norm of the residual r_j, and rgs stands for repeated Gram–Schmidt orthogonalization.

Algorithm 1 Jacobi–Davidson method for the three-parameter eigenvalue problem. In the algorithm, ℓ denotes the size of the subspace after a restart, ε is used in the convergence criterion for an eigenvalue, and $\delta > \varepsilon$ is used to decide whether a Ritz pair is a candidate for TRQI refinement

1: Choose initial matrices $U_j^{(0)} \in \mathbb{C}^{n_j \times \ell}$ with orthonormal columns for $j = 1, 2, 3$.

2: **for** $k = 0, 1, \ldots$ **do**

3:     Compute appropriate Ritz value (σ, τ, ψ) and vector $U_1^{(k)} s_1 \otimes U_2^{(k)} s_2 \otimes U_3^{(k)} s_3$ from the projected three-parameter eigenvalue problem

$$U_1^{(k)^H} A_1 U_1^{(k)} s_1 = \sigma\, U_1^{(k)^H} B_1 U_1^{(k)} s_1 + \tau\, U_1^{(k)^H} C_1 U_1^{(k)} s_1 + \psi\, U_1^{(k)^H} D_1 U_1^{(k)} s_1,$$

$$U_2^{(k)^H} A_2 U_2^{(k)} s_2 = \sigma\, U_2^{(k)^H} B_2 U_2^{(k)} s_2 + \tau\, U_2^{(k)^H} C_2 U_2^{(k)} s_2 + \psi\, U_2^{(k)^H} D_2 U_2^{(k)} s_2,$$

$$U_3^{(k)^H} A_3 U_3^{(k)} s_3 = \sigma\, U_3^{(k)^H} B_3 U_3^{(k)} s_3 + \tau\, U_3^{(k)^H} C_3 U_3^{(k)} s_3 + \psi\, U_3^{(k)^H} D_3 U_3^{(k)} s_3.$$

4:     Compute the residual $r_j = (A_j - \sigma B_j - \tau C_j - \psi D_j)\, u_j$, where $u_j = U_j^{(k)} s_j$, for $j = 1, 2, 3$.

5:     **if** $(\|r_1\|^2 + \|r_2\|^2 + \|r_3\|^2)^{1/2} \leq \delta$ **then**

6:         Refine the Ritz pair by applying $t \geq 0$ steps of the TRQI and update the residuals.

7:         If the refined pair satisfies the selection criterion and $(\|r_1\|^2 + \|r_2\|^2 + \|r_3\|^2)^{1/2} \leq \varepsilon$, then extract the eigenpair and compute the corresponding left eigenvector.

8:     **else**

9:         Solve for $j = 1, 2, 3$ (approximately or exactly) the correction equation

$$\left(I - u_j u_j^H\right)(A_j - \sigma B_j - \tau C_j - \psi D_j)\, v_j = -r_j, \qquad v_j \perp u_j. \tag{17}$$

10:       Expand $U_j^{(k+1)} = \text{rgs}\left(U_j^{(k)}, v_j\right)$ for $j = 1, 2, 3$.

11:       If the dimension of $U_j^{(k+1)}$ is too large, construct new $U_j^{(k+1)} \in \mathbb{C}^{n_j \times \ell}$ for $j = 1, 2, 3$.

12:     **end if**

13: **end for**

In Algorithm 1, we extract one eigenpair at a time. A small projected three-parameter eigenvalue problem is solved in each step. If an eigenpair has converged, then we keep the current subspace, as it may lead to other eigenvalues. Otherwise, we expand the subspace with the addition of a vector that satisfies the correction equation (17) in line 10, where we apply repeated Gram–Schmidt orthogonalization. In what follows, we spell out some of the important details of the algorithm.

Targeting. Depending on the application, a prescribed eigenvalue target can be either a point $(\lambda_0, \mu_0, \eta_0)$ or a plane, for example, $\eta = 0$. For instance, if we take $(0, 0, 0)$ as the target, then we search for eigenvalues with the minimal value of $|\lambda|^2 + |\mu|^2 + |\eta|^2$. In line 3, we select a particular Ritz value (σ, τ, ψ) that is closest to the target and satisfies an additional selection criterion described below.

Selection criterion. The purpose of the selection criterion is to prevent convergence to an eigenvalue that has already been detected. The criterion is based on the following lemma, which is a straightforward generalization of its two-parameter counterpart (see the work of Hochstenbach et al.[6]).

Lemma 1. *Let $(\lambda_1, \mu_1, \eta_1) \neq (\lambda_2, \mu_2, \eta_2)$ be different eigenvalues of the three-parameter eigenvalue problem such that $(\lambda_1, \mu_1, \eta_1)$ is a simple eigenvalue with the right eigenvector $x_1^{(1)} \otimes x_2^{(1)} \otimes x_3^{(1)}$ and the left eigenvector $y_1^{(1)} \otimes y_2^{(1)} \otimes y_3^{(1)}$. If $y_1^{(2)} \otimes y_2^{(2)} \otimes y_3^{(2)}$ is a left eigenvector corresponding to $(\lambda_2, \mu_2, \eta_2)$, then*

(i) $(y_1^{(1)} \otimes y_2^{(1)} \otimes y_3^{(1)})^H \Delta_0 (x_1^{(1)} \otimes x_2^{(1)} \otimes x_3^{(1)}) \neq 0$*, and*

(ii) $(y_1^{(2)} \otimes y_2^{(2)} \otimes y_3^{(2)})^H \Delta_0 (x_1^{(1)} \otimes x_2^{(1)} \otimes x_3^{(1)}) = 0$*.*

Let $(\lambda_q, \mu_q, \eta_q)$ be the eigenvalues that are already extracted along with the corresponding left and right eigenvectors $y_1^{(q)} \otimes y_2^{(q)} \otimes y_3^{(q)}$ and $x_1^{(q)} \otimes x_2^{(q)} \otimes x_3^{(q)}$ for $q = 1, \ldots, m$. The selection criterion outlined next is based on these eigenvectors. In line 3 of Algorithm 1, we select a Ritz value such that the corresponding Ritz vector $u_1 \otimes u_2 \otimes u_3$, with $u_j = U_j^{(k)} s_j$ for

$j = 1, 2, 3$, satisfies

$$\max_{q=1,\ldots,m} \frac{\left|\left(y_1^{(q)} \otimes y_2^{(q)} \otimes y_3^{(q)}\right)^H \Delta_0(u_1 \otimes u_2 \otimes u_3)\right|}{\left|\left(y_1^{(q)} \otimes y_2^{(q)} \otimes y_3^{(q)}\right)^H \Delta_0\left(x_1^{(q)} \otimes x_2^{(q)} \otimes x_3^{(q)}\right)\right|} < \xi \tag{18}$$

for a given $\xi < 1$, for instance, $\xi = 10^{-1}$. Among those Ritz values satisfying the criterion, we choose the one closest to the prescribed target.

Correction equation and preconditioning. When the target is a point $(\lambda_0, \mu_0, \eta_0)$, we solve the correction equation in line 9 approximately by a Krylov subspace method, for example, by the generalized minimal residual method. An important feature of the Jacobi–Davidson method is the preconditioning applied to the correction equation. A good choice for a preconditioner is the inverse of $A_j - \lambda_0 B_j - \mu_0 C_j - \eta_0 D_j$. Since this matrix has size $n_j \times n_j$, where n_j is usually small compared to $n_1 n_2 n_3$, this is a cheap operation.

If the target is the plane $\eta = 0$, then λ_0 and μ_0 are not defined, and we cannot use the preconditioning discussed above. In this case, we often get good results if we solve the correction equation exactly. This is usually feasible, as in many applications, n_j is not large. We employ the expression

$$v_j = -u_j + \left(u_j^H z_j\right)^{-1} z_j$$

for the exact solution of the correction equation (17), where $z_j := (A_j - \sigma B_j - \tau C_j - \psi D_j)^{-1} u_j$ for $j = 1, 2, 3$ (see the work of Sleijpen and Van der Vorst[12] for the details).

Restarts. To keep the computation efficient, we restart Algorithm 1 in line 11 when the subspace becomes too large. As for the choice of the new subspace of dimension ℓ, we employ

$$U_j^{(k+1)} = \text{rgs}\left(u_j^{(k)} + v_j^{(k)}, \ldots, u_j^{(k-\ell+1)} + v_j^{(k-\ell+1)}\right),$$

where $u_1^{(q)} \otimes u_2^{(q)} \otimes u_3^{(q)}$ is the Ritz vector and $v_j^{(q)}$, $j = 1, 2, 3$, is the solution (exact or approximate) of the corresponding correction equation (17) at iteration q. In this way, we build the new search space from the last ℓ eigenvector approximations.

TRQI. The method performs better if we use Jacobi–Davidson up to a point when the residual of a Ritz pair is reasonably small, that is, smaller than δ in line 5, but still not smaller than ε required for a convergence in line 7. Whenever we find such a Ritz pair, we refine it with the TRQI, which is a generalization of the standard Rayleigh quotient iteration and was also applied to a two-parameter eigenvalue problem in the work of Plestenjak.[18]

Next, we provide a brief description of the TRQI. An eigenpair of the three-parameter eigenvalue problem is a zero of the function

$$F(x, y, z, \lambda, \mu, \eta) = \begin{bmatrix} (A_1 - \lambda B_1 - \mu C_1 - \eta D_1)x \\ (A_2 - \lambda B_2 - \mu C_2 - \eta D_2)y \\ (A_3 - \lambda B_3 - \mu C_3 - \eta D_3)z \\ u^H x - 1 \\ v^H y - 1 \\ w^H z - 1 \end{bmatrix},$$

where constant vectors u, v, w, not orthogonal to x, y, z, respectively, are used for normalization. If $(x_k, y_k, z_k, \lambda_k, \mu_k, \eta_k)$ is an approximation for a zero of F, then we may use Newton's method to obtain a new approximation $(x_k + \Delta x_k, y_k + \Delta y_k, z_k + \Delta z_k, \lambda_k + \Delta \lambda_k, \mu_k + \Delta \mu_k, \eta_k + \Delta \eta_k)$. In the TRQI, we start with an eigenvector approximation $x_k \otimes y_k \otimes z_k$, where $\|x_k\| = \|y_k\| = \|z_k\| = 1$. As an approximation $(\lambda_k, \mu_k, \eta_k)$ for the corresponding eigenvalue, we use the tensor Rayleigh quotient

$$\lambda_k = \frac{(x_k \otimes y_k \otimes z_k)^H \Delta_1 (x_k \otimes y_k \otimes z_k)}{(x_k \otimes y_k \otimes z_k)^H \Delta_0 (x_k \otimes y_k \otimes z_k)},$$

$$\mu_k = \frac{(x_k \otimes y_k \otimes z_k)^H \Delta_2 (x_k \otimes y_k \otimes z_k)}{(x_k \otimes y_k \otimes z_k)^H \Delta_0 (x_k \otimes y_k \otimes z_k)},$$

$$\eta_k = \frac{(x_k \otimes y_k \otimes z_k)^H \Delta_3 (x_k \otimes y_k \otimes z_k)}{(x_k \otimes y_k \otimes z_k)^H \Delta_0 (x_k \otimes y_k \otimes z_k)},$$

and set $x_{k+1}, y_{k+1}, z_{k+1}$ equal to the vectors $x_k + \Delta x_k, y_k + \Delta y_k, z_k + \Delta z_k$ that we get from one step of Newton's method with an initial approximation $(x_k, y_k, z_k, \lambda_k, \mu_k, \eta_k)$. In this Newton step, we set $u = x_k$, $v = y_k$, and $w = z_k$.

Note that when none of n_1, n_2, n_3 is large, one step of the TRQI might be less expensive than one iteration of the Jacobi–Davidson method, and it is more efficient to switch to the TRQI to extract the eigenpair once the Jacobi–Davidson method gets close enough. The choice of the parameter δ requires care. If we set δ too large, then the TRQI refinement is applied to poor candidates, and the TRQI might converge to an eigenvalue that is not close to the target or an eigenvalue that is already extracted. On the other hand, if δ is too small, then the condition in line 5 might never be fulfilled, and the method might not return any eigenvalues.

Harmonic Ritz values. Last but not least, let us note that although it is straightforward to generalize harmonic Ritz values from the work of Hochstenbach and Plestenjak[17] to three-parameter eigenvalue problems, we omit this ingredient in the description of the algorithm for simplicity. We do not use harmonic Ritz values in the numerical experiments with the Jacobi–Davidson method in Section 5, but the use of harmonic Ritz values is an option in the implementation of Algorithm 1 in MultiParEig.[19]

4.3 | Use of a Krylov subspace method with full size tensor vectors

To find eigenvalues with the smallest $|\eta|$, we can also consider methods that operate on the generalized eigenvalue problem $\Delta_3 z = \eta \Delta_0 z$. We present some alternatives in this subsection and in the succeeding two subsections.

We consider a Krylov subspace method for (16), which means that in each step, we have to solve a linear system

$$\Delta_3 w = \Delta_0 v \tag{19}$$

for the unknown w efficiently. While we can exploit the connection of such linear systems to Sylvester equations in the two-parameter case, it does not seem straightforward to extend the Sylvester equation approach to the three-parameter setting. Consequently, it remains an open problem how to solve (19) with a complexity below $\mathcal{O}(n_1^3 n_2^3 n_3^3)$.

More specifically, by introducing the vectorizations $v = \text{vec}(\mathcal{V})$ and $w = \text{vec}(\mathcal{W})$, where $\mathcal{V}, \mathcal{W} \in \mathbb{C}^{n_1 \times n_2 \times n_3}$ are three-dimensional tensors, it is possible to express (19) as

$$\mathcal{W} \times_1 B_1 \times_2 C_2 \times_3 A_3 + \mathcal{W} \times_1 C_1 \times_2 A_2 \times_3 B_3 + \cdots - \mathcal{W} \times_1 A_1 \times_2 C_2 \times_3 B_3 = \mathcal{M}, \tag{20}$$

where the right-hand side is $\mathcal{M} = \mathcal{V} \times_1 B_1 \times_2 C_2 \times_3 D_3 + \mathcal{V} \times_1 C_1 \times_2 D_2 \times_3 B_3 + \cdots - \mathcal{V} \times_1 D_1 \times_2 C_2 \times_3 B_3$, and $\times_j$ denotes the j-node product for $j = 1, 2, 3$. Equation (20) resembles a Sylvester equation in three dimensions, but has too many terms. That is, in three dimensions, the Sylvester equation has the form

$$\mathcal{X} \times_1 A + \mathcal{X} \times_2 B + \mathcal{X} \times_3 C = \mathcal{Y}. \tag{21}$$

Using Schur decompositions for matrices A, B, and C, one can solve (21) efficiently by a generalization of the Bartels–Stewart algorithm (see the work of Li et al.[20] for details). Unfortunately, in our setting, we have six nonzero terms in (20), and it does not seem possible to write this equation in the form (21).

4.4 | Subspace iteration

If $n_1 n_2 n_3$ is too large for the approach in the previous subsection, then we can apply subspace iteration to (16) in a way similar to its counterpart for the two-parameter case, using low-rank approximations to make the computation feasible. *The exact subspace iteration* with full vectors operates as follows. We start with a matrix $Z_0 \in \mathbb{C}^{n_1 n_2 n_3 \times p}$ with orthonormal columns. In each step, for a given Z_k, we solve the linear system $\Delta_3 W_k = \Delta_0 Z_k$ for W_k and then set Z_{k+1} equal to the Q factor in the QR decomposition of W_k. Typically, the columns of Z_k converge to an orthonormal basis for the dominant invariant subspace of $\Delta_3^{-1} \Delta_0$, and $Z_k^T \Delta_3^{-1} \Delta_0 Z_k$ converges to an upper triangular matrix with p dominant eigenvalues η of (16) on the diagonal.

As the full columns of Z_k are too large, we use low-rank approximations. We call this variant *inexact subspace iteration*. Specifically, we suppose that all columns of $Z_k \in \mathbb{C}^{n_1 n_2 n_3 \times p}$ lie in a subspace spanned by $U_1^{(k)} \otimes U_2^{(k)} \otimes U_3^{(k)}$ for $U_j^{(k)} \in \mathbb{C}^{n_j \times \ell}$ for $j = 1, 2, 3$. The columns of $Z_k \in \mathbb{C}^{n_1 n_2 n_3 \times p}$ are represented in the Tucker format

$$z_i^{(k)} = \text{vec}\left(\mathcal{D}_i^{(k)} \times_1 U_1^{(k)} \times_2 U_2^{(k)} \times_3 U_3^{(k)} \right), \tag{22}$$

where $\mathcal{D}_i^{(k)}$ is an $\ell \times \ell \times \ell$ core tensor for $i = 1, \ldots, p$, where $p \leq \ell$. Each iteration proceeds as follows.

1. Solve the linear system $\Delta_3 w_i^{(k)} = \Delta_0 z_i^{(k)}$ approximately for $i = 1, \ldots, p$ (see below for details) and orthonormalize the solution vectors $w_1^{(k)}, \ldots, w_p^{(k)}$.

2. Replace the orthonormalized solutions with their low-rank approximations, which leads to $z_1^{(k+1)}, \ldots, z_p^{(k+1)}$ forming the columns of Z_{k+1} for the next step.

In the second step, $z_i^{(k+1)} = \text{vec}(D_i^{(k+1)} \times_1 U_1^{(k+1)} \times_2 U_2^{(k+1)} \times_3 U_3^{(k+1)})$ for some $U_j^{(k+1)} \in \mathbb{C}^{n_j \times \ell}$, $j = 1, 2, 3$. We explain how to form $U_j^{(k+1)}$ in Algorithm 2 at the end of this subsection.

The main part of the inexact subspace iteration is to solve the linear systems $\Delta_3 w_i^{(k)} = \Delta_0 z_i^{(k)}$ for $i = 1, \ldots, p$ approximately by using low-rank approximations. This is justified by the following argument. When v in (19) is an eigenvector of (16), which implies $v = v_1 \otimes v_2 \otimes v_3$ is a decomposable tensor, then the right-hand side of (19) is a sum

$$B_1 v_1 \otimes C_2 v_2 \otimes D_3 v_3 + C_1 v_1 \otimes D_2 v_2 \otimes B_3 v_3 + \cdots - D_1 v_1 \otimes C_2 v_2 \otimes B_3 v_3$$

of six rank-1 tensors. In this case, the solution w of (19) is also an eigenvector of (16) and has rank 1. As in the exact subspace iteration, the columns of Z_k converge to linear combinations of a small number of dominant eigenvectors, and it is reasonable to use low-rank approximations for the solutions of the linear systems $\Delta_3 w_i^{(k)} = \Delta_0 z_i^{(k)}$ for $i = 1, \ldots, p$.

Although we cannot write (20) as a Sylvester equation in the three-parameter setting, we can borrow some ideas from the Krylov method for the two-parameter case that is based on the solutions of Sylvester equations by the low-rank approximation approach due to Hu–Reichel. In particular, suppose that we are looking for a low-rank approximation of the solution of (19). Let us assume that $A_1, A_2,$ and A_3 are nonsingular. Then, (19) is equivalent to

$$\widetilde{\Delta}_3 w := \begin{vmatrix} \widetilde{B}_1 & \widetilde{C}_1 & I \\ \widetilde{B}_2 & \widetilde{C}_2 & I \\ \widetilde{B}_3 & \widetilde{C}_3 & I \end{vmatrix}_\otimes w = \begin{vmatrix} \widetilde{B}_1 & \widetilde{C}_1 & \widetilde{D}_1 \\ \widetilde{B}_2 & \widetilde{C}_2 & \widetilde{D}_2 \\ \widetilde{B}_3 & \widetilde{C}_3 & \widetilde{D}_3 \end{vmatrix}_\otimes v =: \widetilde{\Delta}_0 v,$$

where $\widetilde{\Delta}_3 = (A_1 \otimes A_2 \otimes A_3)^{-1} \Delta_3$, $\widetilde{\Delta}_0 = (A_1 \otimes A_2 \otimes A_3)^{-1} \Delta_0$, $\widetilde{B}_i = A_i^{-1} B_i$, $\widetilde{C}_i = A_i^{-1} C_i$, $\widetilde{D}_i = A_i^{-1} D_i$. Observe that for v of the form $v = \text{vec}(D \times_1 U_1^{(k)} \times_2 U_2^{(k)} \times_3 U_3^{(k)})$ for some D, the vector $\widetilde{\Delta}_0 v$ lies in the subspace spanned by $F_1 \otimes F_2 \otimes F_3$, where $F_j = \text{span}\{\widetilde{B}_j U_j^{(k)}, \widetilde{C}_j U_j^{(k)}, \widetilde{D}_j U_j^{(k)}\}$ for $j = 1, 2, 3$. Our low-rank approach employs the generalized Krylov subspaces

$$\mathcal{K}_r(\widetilde{B}_j, \widetilde{C}_j, F_j) := \text{span}\left\{ M_0(\widetilde{B}_j, \widetilde{C}_j, F_j), M_1(\widetilde{B}_j, \widetilde{C}_j, F_j), \ldots, M_{r-1}(\widetilde{B}_j, \widetilde{C}_j, F_j) \right\} \tag{23}$$

for a modest r, where $M_0(\widetilde{B}_j, \widetilde{C}_j, F_j) = F_j$ and $M_{i+1}(\widetilde{B}_j, \widetilde{C}_j, F_j) = [\widetilde{B}_j M_i(\widetilde{B}_j, \widetilde{C}_j, F_j) \quad \widetilde{C}_j M_i(\widetilde{B}_j, \widetilde{C}_j, F_j)]$ for $i > 0$. This is a generalization of the Krylov subspaces used in the Hu–Reichel method (cf. the work of Weile et al.[21]). An approximate solution of (19) is assumed to be of the form

$$w = \text{vec}(\mathcal{Y} \times_1 Q_1 \times_2 Q_2 \times_3 Q_3) \in \mathcal{K}_r(\widetilde{B}_1, \widetilde{C}_1, F_1) \otimes \mathcal{K}_r(\widetilde{B}_2, \widetilde{C}_2, F_2) \otimes \mathcal{K}_r(\widetilde{B}_3, \widetilde{C}_3, F_3),$$

where Q_j is a matrix whose columns form an orthonormal basis for $\mathcal{K}_r(\widetilde{B}_j, \widetilde{C}_j, F_j)$, and $\mathcal{Y}$ is the solution of the projected equation

$$\begin{vmatrix} Q_1^H \widetilde{B}_1 Q_1 & Q_1^H \widetilde{C}_1 Q_1 & I \\ Q_2^H \widetilde{B}_2 Q_2 & Q_2^H \widetilde{C}_2 Q_2 & I \\ Q_3^H \widetilde{B}_3 Q_3 & Q_3^H \widetilde{C}_3 Q_3 & I \end{vmatrix}_\otimes \text{vec}(\mathcal{Y}) = \begin{vmatrix} Q_1^H \widetilde{B}_1 Q_1 & Q_1^H \widetilde{C}_1 Q_1 & Q_1^H \widetilde{D}_1 Q_1 \\ Q_2^H \widetilde{B}_2 Q_2 & Q_2^H \widetilde{C}_2 Q_2 & Q_2^H \widetilde{D}_2 Q_2 \\ Q_3^H \widetilde{B}_3 Q_3 & Q_3^H \widetilde{C}_3 Q_3 & Q_3^H \widetilde{D}_3 Q_3 \end{vmatrix}_\otimes (Q_1^H \otimes Q_2^H \otimes Q_3^H) v$$

that satisfies the Galerkin condition that the residual is orthogonal to the subspace $\text{span}\{Q_1 \otimes Q_2 \otimes Q_3\}$. In the two-parameter case, we can exploit the relation to the Sylvester equation to solve the projected equation efficiently. As explained in the previous subsection, we are not aware of such a relation in the three-parameter setting. Hence, we solve the projected systems directly. For this reason, the dimension of the subspace $Q_1 \otimes Q_2 \otimes Q_3$ cannot grow too large.

The above procedure yields vectors $w_i^{(k)} = \text{vec}(\mathcal{Y}_i \times_1 Q_1 \times_2 Q_2 \times_3 Q_3)$ for $i = 1, \ldots, p$, which are orthonormalized into $\widetilde{w}_i^{(k)} = \text{vec}(\widetilde{\mathcal{Y}}_i \times_1 Q_1 \times_2 Q_2 \times_3 Q_3)$ for $i = 1, \ldots, p$ by the Gram–Schmidt procedure. We remark that the orthonormalization affects only the core tensors, whereas the subspace bases $Q_1, Q_2,$ and Q_3 do not change. After orthonormalization, we approximate $\widetilde{w}_1^{(k)}, \ldots, \widetilde{w}_p^{(k)}$ by their orthogonal projections onto a low-dimensional subspace $\text{span}\{U_1^{(k+1)} \otimes U_2^{(k+1)} \otimes U_3^{(k+1)}\}$ for some $U_j^{(k+1)} \in \mathbb{C}^{n_j \times \ell}$ such that $\text{span}\{U_j^{(k+1)}\} \subset \text{span}\{Q_j\}$.

Finally, we discuss a feasible approach to construct a suitable $\ell \times \ell \times \ell$-dimensional subspace $\text{span}\{U_1^{(k+1)} \otimes U_2^{(k+1)} \otimes U_3^{(k+1)}\}$ of $\text{span}\{Q_1 \otimes Q_2 \otimes Q_3\}$. As we apply a subspace iteration, we expect that, near convergence, $\widetilde{w}_1^{(k)}$ is close to the dominant eigenvector, which is a decomposable tensor. Furthermore, $\widetilde{w}_2^{(k)}$ should be close to a linear combination of the dominant two eigenvectors, and so on. Thus, we construct the $\ell \times \ell \times \ell$-dimensional subspace

span$\{U_1^{(k+1)} \otimes U_2^{(k+1)} \otimes U_3^{(k+1)}\}$ by considering $\widetilde{w}_1^{(k)}$ first. We determine a subspace span$\{V_1 \otimes V_2 \otimes V_3\}$ that contains a good low-rank approximation of $\widetilde{w}_1^{(k)}$, where V_j is an $n_j \times m_j$ matrix with orthonormal columns and $m_j \leq \ell$ for $j = 1, 2, 3$. While the best low-rank approximation is well defined and easy to compute in the two-parameter case, this is more complicated in the three-parameter setting, where the available tools are the multilinear singular value decomposition (SVD) or a low multilinear rank approximation (see, e.g., the work of Kolda and Bader[22]). Once we obtain a low-rank approximation for $\widetilde{w}_1^{(k)}$, we take V_j as the starting column block of $U_j^{(k+1)}$. Then, we project the next vector $\widetilde{w}_2^{(k)}$ onto the orthogonal complement of span$\{V_1 \otimes V_2 \otimes V_3\}$ and find a new column block for $U_j^{(k+1)}$ from a low-rank approximation of the projected vector. We continue this way until we collect enough columns for $U_j^{(k+1)}$ for $j = 1, 2, 3$. This construction is described in Algorithm 2. Finally, it is worth remarking that for small k, when the subspace is far from an invariant one, we can expect to get all columns of $U_j^{(k+1)}$ for $j = 1, 2, 3$ just from a low-rank approximation of $\widetilde{w}_1^{(k)}$, whereas at later iterations, after $\widetilde{w}_1^{(k)}$ has already converged to a dominant eigenvector, we obtain only the first column of $U_j^{(k+1)}$ from $\widetilde{w}_1^{(k)}$ and the remaining ones from $\widetilde{w}_2^{(k)}, \ldots, \widetilde{w}_p^{(k)}$.

Algorithm 2 Computation of matrices $U_1^{(k+1)}, U_2^{(k+1)}, U_3^{(k+1)}$ with orthonormal columns forming a basis for low-rank approximations of $w_1^{(k)}, \ldots, w_p^{(k)}$, which are the vectors generated by the inexact subspace iteration at step k

---

1: $U_1^{(k+1)} = [], U_2^{(k+1)} = [], U_3^{(k+1)} = [], m = 0$

2: **for** $q = 1, \ldots, p$ and while $m < \ell$ **do**

3: $\quad z = \left(I - \left[U_1^{(k+1)} \otimes U_2^{(k+1)} \otimes U_3^{(k+1)} \right] \left[U_1^{(k+1)} \otimes U_2^{(k+1)} \otimes U_3^{(k+1)} \right]^H \right) w_q^{(k)}$

4: $\quad$ Find matrices V_1, V_2, V_3 with $s \leq \ell - m$ orthonormal columns that are used for a low-rank approximation of z in span$\{V_1 \otimes V_2 \otimes V_3\}$.

5: $\quad U_j^{(k+1)} = \left[U_j^{(k+1)} \ V_j \right]$ for $j = 1, 2, 3$.

6: $\quad m = m + s$.

7: **end for**

4.5 | Subspace iteration with Arnoldi expansion

The inexact subspace iteration for the three-parameter eigenvalue problem presented in the previous subsection is inspired from the ideas in the work of Meerbergen and Plestenjak[8] for the two-parameter case. Here, we further simplify that approach by avoiding the explicit use of low-rank approximations, giving rise to a method that is easier to implement. We will add some new features that are not present in the two-parameter version in the work of Meerbergen and Plestenjak,[8] which improve the efficiency of the approach substantially in the three-parameter case. Some of the new features, for instance, the selection criterion from Section 4.2, can be adopted in the two-parameter version in a straightforward way.

In the inexact subspace iteration of the previous subsection, the approximate solutions of the linear systems at step k are assumed to lie in span$\{Q_1 \otimes Q_2 \otimes Q_3\}$, where the columns of Q_j form an orthonormal basis for the generalized Krylov subspace $\mathcal{K}_r(\widetilde{B}_j, \widetilde{C}_j, F_j)$ defined in (23) for $j = 1, 2, 3$. These spaces contain many approximations for the eigenvectors that we can use to form the next subspace span$\{U_1^{(k+1)} \otimes U_2^{(k+1)} \otimes U_3^{(k+1)}\}$. Here, we form the subspace from ℓ Ritz vectors of the ℓ Ritz values with the smallest $|\psi|$ of the projected three-parameter eigenvalue problem

$$Q_1^H A_1 Q_1 s_1 = \sigma Q_1^H B_1 Q_1 s_1 + \tau Q_1^H C_1 Q_1 s_1 + \psi Q_1^H D_1 Q_1 s_1$$
$$Q_2^H A_2 Q_2 s_2 = \sigma Q_2^H B_2 Q_2 s_2 + \tau Q_2^H C_2 Q_2 s_2 + \psi Q_2^H D_2 Q_2 s_2 \qquad (24)$$
$$Q_3^H A_3 Q_3 s_3 = \sigma Q_3^H B_3 Q_3 s_3 + \tau Q_3^H C_3 Q_3 s_3 + \psi Q_3^H D_3 Q_3 s_3.$$

As each Ritz vector is decomposable, we form $U_j^{(k+1)}$ such that its columns form an orthonormal basis for the subspace spanned by the j-nodes of the selected Ritz vectors. Note that this approach is close to the methods based on tensor decompositions such as those in the works of Kressner et al.[23] and Kressner and Tobler.[24] The main difference is that our approach only uses the factor matrices Q_1, Q_2, Q_3 of a Tucker decomposition, that is, the core tensor is not used. A formal description of the approach is given in Algorithm 3, and some details are discussed below.

Algorithm 3 Subspace iteration with Arnoldi expansion and restarts based on selected Ritz vectors for the generalized eigenvalue problem (19) associated with the three-parameter eigenvalue problem

In the algorithm, ℓ denotes the size of the subspace after a restart, r is the number of block Arnoldi steps, ε is used in the convergence criterion for an eigenvalue, and $\delta > \varepsilon$ controls when a Ritz pair is a candidate for the TRQI refinement

---

1: Choose initial matrices $U_j^{(0)} \in \mathbb{C}^{n_j \times \ell}$ with orthonormal columns for $j = 1, 2, 3$.
2: **for** $k = 0, 1, \ldots$ **do**
3: 　　**for** $j = 1, 2, 3$ **do**
4: 　　　　$F_j = \left[A_j^{-1} B_j U_j^{(k)} \ \ A_j^{-1} C_j U_j^{(k)} \ \ A_j^{-1} D_j U_j^{(k)} \right]$
5: 　　　　Form Q_j whose columns are orthonormal basis for $\mathcal{K}_r \left(A_j^{-1} B_j, A_j^{-1} C_j, F_j \right)$ using a block Arnoldi algorithm with SVD filtering; see Algorithm 4.
6: 　　**end for**
7: 　　**if** the size of $Q_1 \otimes Q_2 \otimes Q_3$ is too large **then**
8: 　　　　Shrink matrices Q_1, Q_2, Q_3 by the same factor by removing the appropriate number of the last columns.
9: 　　**end if**
10: 　　Compute m Ritz values $(\sigma_i, \tau_i, \psi_i)$ and Ritz vectors $z_1^{(i)} \otimes z_2^{(i)} \otimes z_3^{(i)} := Q_1 s_1^{(i)} \otimes Q_2 s_2^{(i)} \otimes Q_3 s_3^{(i)}$ for $i = 1, \ldots, m$ with the smallest values of $|\psi|$ from the projected three-parameter eigenvalue problem (24).
11: 　　Refine Ritz pairs $\left((\sigma_i, \tau_i, \psi_i), z_1^{(i)} \otimes z_2^{(i)} \otimes z_3^{(i)} \right)$ for $i = 1, \ldots, m$ by $s \geq 0$ steps of the TRQI.
12: 　　**for** $i = 1, \ldots, m$ **do**
13: 　　　　**if** the Ritz pair $\left((\sigma_i, \tau_i, \psi_i), z_1^{(i)} \otimes z_2^{(i)} \otimes z_3^{(i)} \right)$ satisfies the selection criterion **then**
14: 　　　　　　Compute the residual $r_{ij} = (A_j - \sigma_i B_j - \tau_i C_j - \psi_i D_j) z_j^{(i)}$ for $j = 1, 2, 3$.
15: 　　　　　　**if** $(\|r_{i1}\|^2 + \|r_{i2}\|^2 + \|r_{i3}\|^2)^{1/2} \leq \delta$ **then**
16: 　　　　　　　　Further refine the Ritz pair with $t \geq 0$ steps of the TRQI and update the residuals.
17: 　　　　　　　　**if** refined pair satisfies the selection criterion **and** $(\|r_{i1}\|^2 + \|r_{i2}\|^2 + \|r_{i3}\|^2)^{1/2} \leq \varepsilon$ **then**
18: 　　　　　　　　　　Extract the eigenpair and compute the corresponding left eigenvector
19: 　　　　　　　　**end if**
20: 　　　　　　**end if**
21: 　　　　**end if**
22: 　　**end for**
23: 　　Let $p_1, \ldots, p_\ell$ be the indices of the first ℓ Ritz pairs that satisfied the selection criterion, but did not lead to an eigenpair.
24: 　　Form $U_j^{(k+1)}$ whose columns make an orthonormal basis for span $\left\{ z_j^{(p_1)}, \ldots, z_j^{(p_\ell)} \right\}$ for $j = 1, 2, 3$.
25: **end for**

Block Arnoldi algorithm with SVD filtering. The block Arnoldi algorithm in line 5 employed together with an SVD filtering is presented in Algorithm 4. In the three-parameter setting, we are quite limited in the maximum search space. In particular, if the size of the subspace span$\{Q_1 \otimes Q_2 \otimes Q_3\}$ is too large, then we cannot solve the projected problem in line 10. Hence, we use the SVD filtering and the relative cutoff parameter $\zeta \geq 0$ to prevent, on the one hand, the search space from growing too much and, on the other hand, to keep all the significant directions in the subspace. In our experiments, $\zeta = 10^{-5}$ gives good results in practice.

Selection criterion. In line 13 of Algorithm 3, we use the same selection criterion as in the Jacobi–Davidson method, defined by (18). As we need the left eigenvectors corresponding to the eigenvalues that are already extracted to check this criterion, we compute a left eigenvector in line 18 for each new eigenvalue that we find. If a Ritz pair satisfies the selection criterion, it can still happen that the TRQI refinement converges to one of the eigenvalues already extracted. Therefore, we test the selection criterion in line 17 once again to make sure that an eigenvalue is not repeated.

TRQI refinement. The convergence can be drastically improved if we refine all Ritz pairs with a small number of TRQI steps in line 11 of Algorithm 3. This improves the directions that we use for a restart in line 24; additionally, it yields more candidates that satisfy the criterion in line 15. However, we should not use too many refinement steps because even when the TRQI is applied to a poor approximation, it can still converge to an eigenpair. In most cases, such a converged eigenpair is not close to the prescribed target (e.g., it does not have a small $|\eta|$) or is an eigenpair that is already extracted.

If, after this initial TRQI refinement, the selection criterion is satisfied by a Ritz pair and the norm of the corresponding residual is below δ in line 15, then the TRQI refinement is applied once again to the candidate Ritz pair. As in Algorithm 1, the parameter δ should be chosen with care. Since we use only a few steps of block Arnoldi to form our search space, we cannot expect it to contain very good approximations of the eigenvectors. Hence, we perform the second stage of the TRQI on approximations with residuals that are reasonably small to overcome their inaccuracy due to the crudeness of the subspaces.

Algorithm 4 Block Arnoldi expansion with an SVD filtering to form an orthonormal basis for $\mathcal{K}_r(B, C, F)$. In the algorithm, $\zeta \geq 0$ denotes the relative cutoff parameter for the singular values

1: Compute the singular value decomposition $F = U\Sigma V^T$.
2: Select $W = [u_1 \cdots u_j]$, where j is such that $\sigma_j \geq \zeta\sigma_1 > \sigma_{j+1}$, or j is the number of columns of F.
3: $Q = W$
4: **for** $k = 1, \ldots, r$ **do**
5:     $G = (I - QQ^H)[BW \quad CW]$
6:     Compute the singular value decomposition $G = U\Sigma V^T$.
7:     Select $W = [u_1 \cdots u_j]$, where j is such that $\sigma_j \geq \zeta\sigma_1 > \sigma_{j+1}$, or j is the number of columns of G.
8:     $Q = [Q \quad W]$
9: **end for**

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5 | NUMERICAL RESULTS

Algorithms 1 and 3 are both implemented in the MATLAB package MultiParEig.[19] In this section, we conduct numerical experiments with these implementations on several three-parameter eigenvalue problems; all of these examples are available in MultiParEig. The results have been obtained using MATLAB R2012b on a PC having 16-GB RAM and an i5-4670 3.4-GHz CPU.

5.1 | Ellipsoidal wave equation

The first two numerical experiments are performed on the ellipsoidal wave equation described in Section 2.1 with the particular choices $x_0 = 1, y_0 = 1.5$, and $z_0 = 2$ for the radii of the semi-axes of the ellipsoid and $\rho = \sigma = \tau = 0$ for the configuration. This problem was solved numerically using matrices of size 25×25 and the approach from Section 4.3 in the work of Plestenjak et al.[5] Using Algorithms 1 and 3, we can work with much larger matrices corresponding to finer discretizations and obtain more accurate results for the low eigenfrequencies.

We discretize (6) using the Chebyshev collocation on 300 points. We know that all eigenvalues (λ, μ, η) of (6) are real and such that $\eta > 0$ (see, e.g., the work of Levitina[10]). As we are interested in eigenvalues with η closest to the target $\eta_{\text{tar}} \geq 0$, we apply the substitution $(\widetilde{\lambda}, \widetilde{\mu}, \widetilde{\eta}) = (\lambda + 5, \mu, \eta - \eta_{\text{tar}})$ and search for eigenvalues close to $\widetilde{\eta} = 0$ of the transformed problem, with the coefficient matrices $\widetilde{A}_j = A_j + 5B_j - \eta_{\text{tar}}D_j$, $\widetilde{B}_j = B_j$, $\widetilde{C}_j = C_j$, and $\widetilde{D}_j = D_j$ for $j = 1, 2, 3$. We use the shift 5 (where 5 is more or less randomly chosen) to make $\widetilde{A}_j$ nonsingular in the case $\eta_{\text{tar}} = 0$; it changes the λ components of the eigenvalues but does not affect our search, which is based on a prescribed target on the η components of the eigenvalues.

Before applying the numerical methods, we multiply the jth equation by $\widetilde{A}_j^{-1}$ for $j = 1, 2, 3$, after ensuring that $\widetilde{A}_j$ is nonsingular. This is equivalent to considering the generalized eigenvalue problem

$$(\widetilde{A}_1 \otimes \widetilde{A}_2 \otimes \widetilde{A}_3)^{-1}\widetilde{\Delta}_3 z = \eta(\widetilde{A}_1 \otimes \widetilde{A}_2 \otimes \widetilde{A}_3)^{-1}\widetilde{\Delta}_0 z \tag{25}$$

instead of $\widetilde{\Delta}_3 z = \eta\widetilde{\Delta}_0 z$. We do this because the Chebyshev collocation returns matrices such that $\|\widetilde{A}_j\| \gg \|\widetilde{B}_j\|, \|\widetilde{C}_j\|, \|\widetilde{D}_j\|$, and $\widetilde{A}_j$ is ill-conditioned for $j = 1, 2, 3$, where $\|\cdot\|$ denotes the matrix 2-norm. These facts, in turn, imply that $\|\widetilde{\Delta}_3\| \gg \|\widetilde{\Delta}_0\|$, and $\widetilde{\Delta}_3$ is ill-conditioned. We expect that Ritz values of (25) are better approximations for the eigenvalues with the smallest value of $|\eta|$.

Example 1 (Jacobi–Davidson on the ellipsoidal wave equation).
We apply Algorithm 1, where we set the plane $\eta = 0$ as the target and solve the correction equation exactly. We restrict the subspace dimensions between 5 and 10; in particular, we restart using the eigenvector approximations from the

last five iterations. In line 3, the Ritz values are arranged in increasing order according to their distances from the target. We consider a Ritz pair as a candidate for an eigenpair if its residual is smaller than $\delta = 10^{-1}$ and if it satisfies the selection criterion (18) with $\xi_1 = 10^{-1}$. In this case, we refine the Ritz pair with up to four steps of the TRQI. After the refinement, if the residual drops below $\varepsilon = 10^{-8}$ and if the selection criterion with $\xi_2 = 10^{-4}$ is satisfied, then the refined pair is accepted as a new eigenpair. We can extract more than one eigenvalue from the same subspace (without executing the "else" part of the "if" statement, that is, without executing lines 9–11).

We have computed 80 eigenvalues for the following three cases:

a) $\eta_{tar} = 0$, the target eigenvalues are exterior;
b) $\eta_{tar} = 200$, the desired eigenvalues are close to the exterior ones, since there are only a few hundred eigenvalues with their η component satisfying $\eta < 200$;
c) $\eta_{tar} = 1,000$, the target eigenvalues are mildly interior and more difficult to compute.

The computational times for cases a), b), and c) are 7, 480, and 540 seconds for 15, 472, and 594 iterations, respectively. Figure 2 shows the values of $|\eta - \eta_{tar}|$ of the computed eigenvalues in the order of retrieval. In case a), the eigenvalues converge almost in the desired order. In case b), the eigenvalues are not computed in such a desirable order (i.e., the monotonicity of the distances of the η components to the prescribed target with respect to the order of the retrieval degrades slightly), but the method still extracts the eigenvalues close to the prescribed target. In case c), the eigenvalues are retrieved even in a less-structured order, and we need to compute many eigenvalues to be sure that we get the desired eigenvalues closest to the target.

We explored how many eigenvalues one needs to compute with the above settings to get the first 40, 20, and 10 eigenvalues with their η components closest to η_{tar} for cases a), b), and c), respectively. We decrease the number of targeted eigenvalues for larger values of η_{tar}, as interior eigenvalues are more difficult to compute. In Table 1, we report the average results together with the best and the worst run of the algorithm over a set of 10 different random initial subspaces. To make sure that we have all of the closest eigenvalues so that the comparisons are fair, we have computed the eigenvalues a priori repeatedly several times.

In Table 2, we provide the three computed eigenvalues closest to the target for cases a), b), and c). While the closest three eigenvalues for case a) have already been listed in the work of Plestenjak et al.,[5] it was not possible then to compute accurate solutions for cases b) and c) as this requires matrices larger than the methods at that time could handle.
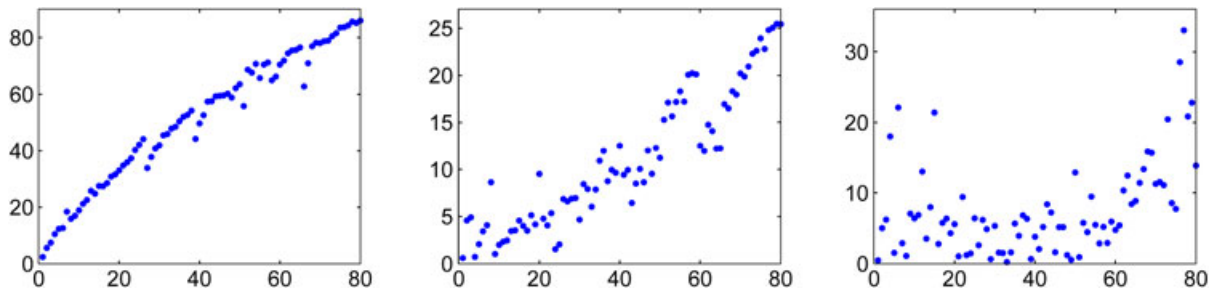

FIGURE 2 Jacobi–Davidson method for the three-parameter eigenvalue problem in Section 5.1. The values $|\eta - \eta_{tar}|$ (vertical axis) of the first 80 computed eigenvalues (λ, μ, η) are plotted with respect to the order of retrieval (horizontal axis) for cases a) $\eta_{tar} = 0$ (left), b) $\eta_{tar} = 200$ (middle), and c) $\eta_{tar} = 1,000$ (right)

TABLE 1 The Jacobi–Davidson method for the three-parameter eigenvalue problem in Section 5.1. The total number of eigenvalues that had to be computed, the number of subspace updates that had to be performed, and computational times in order to retrieve the targeted number of eigenvalues (λ, μ, η) with their η components closest to η_{tar} are listed

η_{tar}	targeted no.	No. of computed eigenvalues			No. of subspace updates			Time (s)		
		average	min	max	average	min	max	average	min	max
0	40	40	40	40	9.3	9	10	3.9	2.7	7.1
200	20	33.8	24	42	199.8	121	311	159.9	64.5	285.1
1,000	10	118.1	51	170	871.2	292	1354	771.3	264.8	1246.5

TABLE 2 A list of three eigenvalues (λ, μ, η) with their η components closest to the targets $\eta_{\text{tar}} = 0, \eta_{\text{tar}} = 200$, and $\eta_{\text{tar}} = 1,000$ for the three-parameter eigenvalue problem resulting from the ellipsoidal wave equation (6) with the radii values $x_0 = 1, y_0 = 1.5$, and $z_0 = 2$ and for the configuration $(\rho, \sigma, \tau) = (0, 0, 0)$. The eigenfrequencies ω corresponding to these computed eigenvalues are also listed in the last column

Target η_{tar}	λ	μ	η	ω
	0.84989209	−3.75231782	2.40498182	2.34458979
0	7.22643744	−13.03122756	5.59866649	3.57728277
	2.05458475	−13.46994828	7.46473320	4.13064732
	141.38925861	−404.17476271	200.60583308	21.41325801
200	63.08832970	−423.06537129	199.27518005	21.34212093
	317.06224687	−551.46171960	201.03180983	21.43598096
	1413.79140334	−2535.12357474	999.75548115	47.80329890
1,000	366.26819031	−2143.09786044	1000.47359673	47.82046416
	1725.45584215	−2758.97471801	999.44259731	47.79581804

Example 2 (Subspace iteration on the ellipsoidal wave equation)**.**
We aim to compute the same eigenvalues as in the previous example using Algorithm 3. Initially, we choose the search space of dimension $\ell = 6$ and apply zero Arnoldi steps (in cases a) and b)) or one Arnoldi step (in case c)) in the expansion. We apply SVD filtering to F_j as in line 4, where we set the cutoff parameter $\zeta = 10^{-5}$. The dimension of $\text{span}(Q_1 \otimes Q_2 \otimes Q_3)$ in line 8 is limited to 1,000, 5,000, and 15,000 in cases a), b), and c), respectively. Smaller subspace dimensions are sufficient when eigenvalues with smaller η components are targeted as these eigenvalues lie in the exterior of the spectrum. On the other hand, larger subspaces are needed for larger values of η_{tar}. In every iteration, we compute 100 (in cases a) and b)) or 50 (in case c)) Ritz values of the projected three-parameter eigenvalue problem in line 10 closest to the prescribed target. This is followed by one step (in cases a) and b)) or three steps (in case c)) of the TRQI to refine each Ritz pair in line 11. After that, we consider a Ritz pair as a candidate for an eigenpair if its residual is smaller than $\delta = 10^{-2}$ and if it satisfies the selection criterion (18) with $\xi_1 = 10^{-1}$. In this case, we refine the Ritz pair with up to three additional steps of the TRQI. The final residuals corresponding to the Ritz pairs are accepted small enough with the particular choices of the parameters as in the Jacobi–Davidson method, that is, $\varepsilon = 10^{-8}$ and $\xi_2 = 10^{-4}$ in line 18.

We computed 80 eigenvalues for cases a), b), and c) from Example 1. The computational times for cases a), b), and c) are 8, 60, and 421 s, and three, three, and five subspace iterations have been carried out, respectively. Figure 3 shows the values of $|\eta - \eta_{\text{tar}}|$ for the computed eigenvalues with respect to their order of retrieval; one can observe a behavior similar to the Jacobi–Davidson method, that is, for smaller values of η_{tar}, it is possible to observe a monotonicity in $|\eta - \eta_{\text{tar}}|$ relative to the order of the retrieval of the eigenvalues, which gradually degrades as η_{tar} is increased.

The SVD filtering does not reduce the dimension of the subspaces enough; hence, we also have to perform the shrinking in line 8 of Algorithm 3. A stricter SVD filtering with a larger cutoff is not a solution, as this results in the removal of some of the good search spaces. For case b), the dimensions of $\text{span}(Q_1 \otimes Q_2 \otimes Q_3)$ are 154,548, 48,300, and 11,340 in the first, second, and third iterations, respectively, which are all shrunk into subspaces of dimension smaller than 5,000. The appearance of larger subspaces in the initial iterations is typical. In the first iteration, the Arnoldi expansion increases the dimension of the search space considerably, but, after a few subspace iterations, the search space contains good approximations of the eigenvectors, and the Arnoldi expansion does not yield many independent directions. These findings are in line with results obtained for the two-parameter case.[25]

Following the practice in Example 1, we explored how many eigenvalues need to be computed in total with the above settings in order to retrieve all of the 40, 20, and 10 eigenvalues with η components closest to η_{tar} for cases a), b), and c), respectively. The results are reported in Table 3.

When we compare the numerical results obtained for the Jacobi–Davidson method and the subspace iteration method, we see that the subspace iteration works slightly faster for mildly interior eigenvalues. This comes at the expense of much larger memory requirements; for instance, at least 16 GB of RAM is needed by subspace iteration to use a search space of dimension 15,000. If subspaces are restricted to small dimensions, then we do not get approximations that are good enough to lead to eigenpairs (even if additional subspace iterations are allowed).
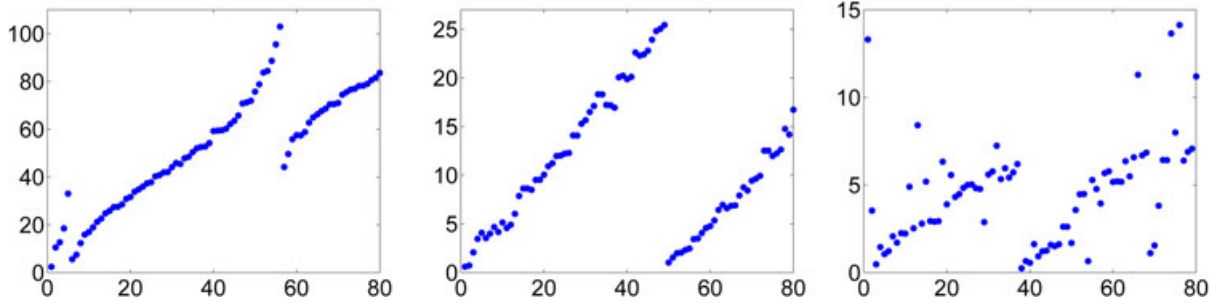
FIGURE 3 Application of the subspace iteration with Arnoldi expansion to the three-parameter eigenvalue problem in Section 5.1. The values $|\eta - \eta_{tar}|$ (vertical axis) of the first 80 computed eigenvalues (λ, μ, η) are plotted with respect to their order of retrieval (horizontal axis) for cases a) $\eta_{tar} = 0$ (left), b) $\eta_{tar} = 200$ (middle), and c) $\eta_{tar} = 1,000$ (right)

TABLE 3 This Table concerns the application of subspace iteration with Arnoldi expansion to the three-parameter eigenvalue problem in Section 5.1. The total number of eigenvalues that had to be computed, subspace iterations that had to be performed, and the computational times in order to retrieve the targeted number of eigenvalues (λ, μ, η) with η components closest to η_{tar} are listed

η_{tar}	targeted no.	No. of computed eigenvalues			No. of subspace iterations			Time (s)		
		average	min	max	average	min	max	average	min	max
0	40	58.2	40	82	2.7	2	3	7.7	5.5	10.3
200	20	74.1	40	109	3.5	3	4	59.8	50.8	69.2
1,000	10	68.9	50	88	5.3	4	7	461.2	307.4	651.7

5.2 | Baer wave equations

By solving the three-parameter eigenvalue problem resulting from the Baer wave equations discussed in Section 2.2, we can obtain many estimates for low eigenfrequencies of the Helmholtz equation (5) on the specified intersection of paraboloids. We could not find any similar numerical results regarding this example in the literature; thus, up to our knowledge, this is the first time that the Helmholtz equation is solved numerically in paraboloidal coordinates. The results could be used for future comparisons to other numerical methods.

As in the previous subsection, we discretize the system of the Baer wave equation (9) for the configuration $(\rho, \sigma) = (0, 0)$ with Chebyshev collocation on 300 points. We are interested in

a) the lowest eigenfrequencies (i.e., $\eta_{tar} = 0$) and
b) the eigenfrequencies closest to 10 (i.e., $\eta_{tar} = 100$).

In case b), we apply the substitution $(\tilde{\lambda}, \tilde{\mu}, \tilde{\eta}) = (\lambda, \mu, \eta - \eta_{tar})$ and search for eigenvalues close to $\tilde{\eta} = 0$ of the transformed problem, with the coefficient matrices $\tilde{A}_j = A_j - \eta_{tar} D_j$, $\tilde{B}_j = B_j$, $\tilde{C}_j = C_j$, and $\tilde{D}_j = D_j$ for $j = 1, 2, 3$. Once again, before applying the numerical methods, we multiply the jth equation by $\tilde{A}_j^{-1}$ for $j = 1, 2, 3$.

Example 3 (Results for Baer wave equations).
We apply both algorithms to the problem above. Using the same settings as in Example 1, the Jacobi–Davidson method computes 80 eigenvalues in 10 s after 16 subspace updates in case a) and in 314 s using 341 subspace updates in case b). For subspace iteration, we use the same settings as in cases a) and b) of Example 2. This means that we limit the dimension of the search space to 1,000 in case a) and to 5,000 in case b). The method requires 7 s and two subspace iterations to compute 80 eigenvalues in case a) and 59 s and three subspace iterations in case b).

We omit the plots of $|\eta - \eta_{tar}|$ with respect to the retrieval order for the converged eigenvalues, as they turn out to be similar to the left-hand and middle plots in Figures 2 and 3. As in Examples 1 and 2, we end up computing more eigenvalues for larger values of η_{tar} in order to retrieve all of the desired eigenvalues closest to η_{tar}.

Similar to the previous examples, we tested how many eigenvalues need to be computed with the settings above in order to retrieve all of the 40 and 20 eigenvalues with η components closest to η_{tar} for cases a) and b), respectively. For both methods, Table 4 reports the average results together with the best and the worst run over a set of 10 different random initial subspaces.

TABLE 4 The Jacobi–Davidson method (JD) and the subspace iteration with Arnoldi expansion (SI) applied to the three-parameter eigenvalue problem in Section 5.2. This Table lists the total number of eigenvalues that had to be computed, the number of subspace iterations that had to be performed, and computational times required to retrieve all of the targeted eigenvalues (λ, μ, η) with their η components closest to η_{tar}

method	η_{tar}	targeted no.	No. of computed eigenvalues			No. of subspace iterations			Time (s)		
			average	min	max	average	min	max	average	min	max
JD	0	40	40	40	40	8.8	7	10	3.4	1.8	7.6
JD	100	20	70.6	41	142	268.5	137	453	209.7	104.0	356.0
SI	0	40	44.0	40	80	2.1	2	3	5.9	5.4	9.4
SI	100	20	88.6	63	104	4.3	4	5	74.2	68.0	86.2

TABLE 5 Results for the Helmholtz equation with a Dirichlet boundary condition on a domain bounded by two elliptic paraboloids $\gamma = 0$ and $\beta = 5$ in paraboloidal coordinates with $c = 1$ and $b = 3$ for the configuration $(\sigma, \rho) = (0, 0)$. Estimates for the lowest six eigenfrequencies and the first three eigenfrequencies larger than 10 of a related three-parameter eigenvalue problem, namely, the Baer wave equation (7), are listed in the Table. In each row, in addition to the eigenfrequency ω, the corresponding eigenvalue (λ, μ, η) and the indices (j_1, j_2, j_3) of the corresponding functions X_1, X_2, X_3 are also listed

λ	μ	η	ω	j_1	j_2	j_3
4.68572309	−4.68336498	1.06171767	1.03039685	0	0	0
8.98735825	−10.98752097	2.52640136	1.58946575	0	1	0
7.84880354	−9.81384367	2.70641882	1.64511970	0	0	1
23.88802753	−18.11389297	3.33102584	1.82510982	1	0	0
15.35149716	−20.44266626	4.60326049	2.14552103	0	2	0
13.98083910	−19.03124115	4.90993954	2.21583834	0	1	1
⋮	⋮	⋮	⋮	⋮	⋮	⋮
368.61672638	−467.93904610	100.12807872	10.00640189	3	10	2
909.43143081	−643.56267025	100.20818157	10.01040367	9	4	0
315.21740925	−436.37381658	100.32096431	10.01603536	2	10	3

Algorithms 1 and 3 return the same 10 eigenfrequencies closest to 10. In particular, the results by both algorithms agree on the first three eigenfrequencies larger than 10; these eigenfrequencies are listed in Table 5 along with the lowest six eigenfrequencies from case a). We verify the correctness of the computed results by means of the Klein oscillation property, which concerns the number of zeros of $X_i(\xi_i)$ as in (7). This property is formally stated in the next theorem. To our knowledge, it has not been explicitly shown for the system of Baer wave equations up to this point; hence, a proof is included in the Appendix.

Theorem 1. *For each of the four possible configurations (σ, τ) in (8), the system of Baer wave differential equations (7) has the Klein oscillation property, that is, all of its eigenvalues are real, and for each triple of nonnegative integers (j_1, j_2, j_3), there exists exactly one eigenvalue (λ, μ, η) such that the corresponding eigenfunctions $X_1(\xi_1)$, $X_2(\xi_2)$, and $X_3(\xi_3)$ have exactly j_1 zeros on (γ, c), j_2 zeros on (c, b), and j_3 zeros on (b, β), respectively.*

In Table 5, we provide an integer triple (j_1, j_2, j_3) for each eigenfrequency ω, with j_i denoting the index of $X_i(\xi_i)$ as in (7), that is, the number of the zeros of the corresponding solution $X_i(\xi_i)$ on the interval (ℓ_i, ℓ_{i+1}) with $\ell_1 = \gamma = 0$, $\ell_2 = c = 1$, $\ell_3 = b = 3$, and $\ell_4 = \beta = 5$. The reported results in the Table are in harmony with Theorem 1, that is, there exists exactly one eigenvalue corresponding to each nonnegative triple (j_1, j_2, j_3). Furthermore, the results confirm that the lowest eigenfrequencies have the smallest indices, as expected in theory (see section 8 in the work of Atkinson and Mingarelli[4]).

Note that we can approximate the solutions $X_1(\xi_1)$, $X_2(\xi_2)$, and $X_3(\xi_3)$ of the Baer wave equation by employing (9) subject to the boundary conditions (10), as well as eigenvectors of the discretized algebraic three-parameter eigenvalue problem. We can combine them in a smooth eigenfunction $X(\xi)$ bounded at the points $\xi = 1, \xi = 3$ and satisfying

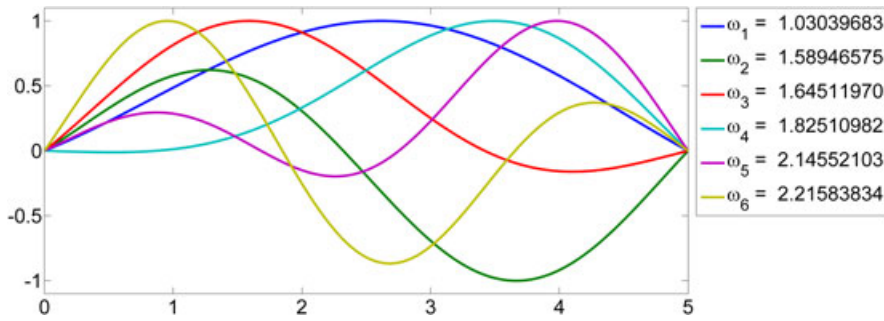$$(\xi - 1)(\xi - 3)X'' + \frac{1}{2}(2\xi - 4)X' + (\lambda + \mu\xi + \eta\xi^2)X = 0$$

FIGURE 4 The eigenfunctions corresponding to the six lowest eigenfrequencies in Table 5

over $\xi \in [0, 5]$ subject to $X(0) = 0$, $X(5) = 0$. The eigenfunctions corresponding to the six lowest eigenfrequencies computed are displayed in Figure 4.

5.3 | Randomly generated example

Our final example is a three-parameter eigenvalue problem generated in MATLAB in such a way that we know all of the eigenvalues. We first form the matrices

$$A_i = U_i \operatorname{diag}(a_i) V_i, \quad B_i = U_i \operatorname{diag}(b_i) V_i, \quad C_i = U_i \operatorname{diag}(c_i) V_i, \quad D_i = U_i \operatorname{diag}(d_i) V_i,$$

where a_i, b_i, c_i, d_i are real random vectors of size n, and U_i, V_i are random well-conditioned sparse matrices of size $n \times n$ for $i = 1, 2, 3$. Observe that the eigenvalues of the resulting three-parameter eigenvalue problem are the solutions to the 3×3 linear systems

$$(a_1)_\ell = \lambda (b_1)_\ell + \mu (c_1)_\ell + \eta (d_1)_\ell,$$
$$(a_2)_j = \lambda (b_2)_j + \mu (c_2)_j + \eta (d_2)_j,$$
$$(a_3)_k = \lambda (b_3)_k + \mu (c_3)_k + \eta (d_3)_k$$

for $\ell, j, k = 1, \ldots, n$, where $(a_i)_p, (b_i)_p, (c_i)_p, (d_i)_p$ denote the pth entries of a_i, b_i, c_i, d_i for $i = 1, 2, 3$. This observation enables us to compute all n^3 eigenvalues for moderate values of n, for example, $n = 100$.

Example 4 (Randomly generated example).
We set $n = 100$ and generate entries of a_i, b_i, c_i, d_i randomly, by first selecting them independently from a uniform distribution over $[0, 1]$ and then applying shifts. The MATLAB code generating a_i, b_i, c_i, d_i and the matrices U_i, V_i for $i = 1, 2, 3$ is given below.

```
U1 = 0.3*sprand(n,n,0.04)+speye(n); U2 = 0.3*sprand(n,n,0.04)+speye(n);
U3 = 0.3*sprand(n,n,0.04)+speye(n); V1 = 0.3*sprand(n,n,0.04)+speye(n);
V2 = 0.3*sprand(n,n,0.04)+speye(n); V3 = 0.3*sprand(n,n,0.04)+speye(n);
a1 = rand(n,1)-0.5; b1 = rand(n,1)+2; c1 = rand(n,1); d1 = rand(n,1)-1;
a2 = rand(n,1)-0.5; b2 = rand(n,1); c2 = rand(n,1)+2; d2 = rand(n,1)+0.5;
a3 = rand(n,1)-0.5; b3 = rand(n,1)-1; c3 = rand(n,1); d3 = rand(n,1)+2;
```

Figure 5 illustrates the resulting eigenvalues projected orthogonally onto the plane $\mu = 0$. Orthogonal projections of the eigenvalues onto the planes $\lambda = 0$ and $\eta = 0$ yield similar pictures.
 We test the methods to compute

a) 20 external eigenvalues with η closest to $\eta_{\text{tar}} = -0.8$ and
b) 10 mildly interior eigenvalues with η closest to $\eta_{\text{tar}} = -0.5$.

Following the practice in the other examples, we use the substitution $(\widetilde{\lambda}, \widetilde{\mu}, \widetilde{\eta}) = (\lambda, \mu, \eta - \eta_{\text{tar}})$ and search for the eigenvalues of the transformed problem having $|\widetilde{\eta}|$ as small as possible.
 We apply the Jacobi–Davidson method, where we solve the correction equation exactly, and use up to three TRQI steps as well as the choices $\delta = 10^{-6}$ and $\varepsilon = 10^{-10}$ to decide whether the residual of a Ritz pair is small enough to
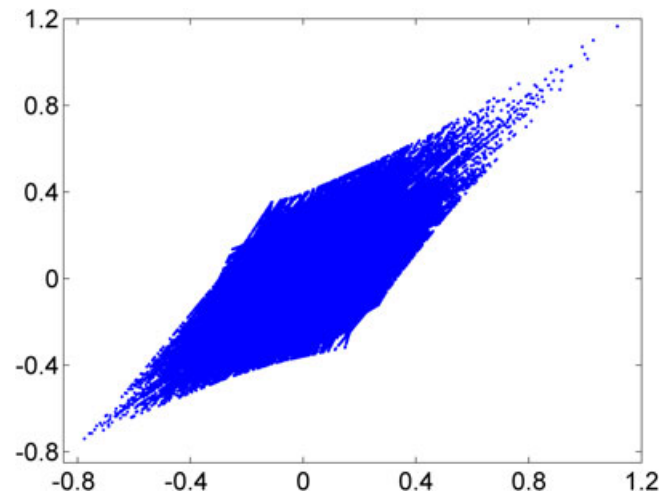
FIGURE 5 The orthogonal projections of the eigenvalues of the three-parameter eigenvalue problem considered in Example 4 onto the plane $\mu = 0$. The horizontal and vertical axes correspond to the λ and η components, respectively

TABLE 6 The Jacobi–Davidson method applied to a random three-parameter eigenvalue problem in Example 4. The Table reports the number of eigenvalues that had to be computed, subspace iterations that had to be performed, and computational times required in order to retrieve all of the targeted eigenvalues (λ, μ, η) with η components closest to η_{tar}

		No. of computed eigenvalues			No. of subspace updates			Time (s)		
η_{tar}	targeted no.	average	min	max	average	min	max	average	min	max
−0.8	20	78.5	38	137	206.3	112	363	197	102	346
−0.5	10	97.3	45	195	284.4	148	478	266	135	454

consider it as an eigenpair. All of the remaining parameters are as in the previous examples. The results are presented in Table 6.

The subspace iteration does not work well on this example. We could not find a combination of parameters to make it competitive with the Jacobi–Davidson method. The method computes some eigenpairs but requires a lot of time and returns many eigenvalues far away from the target.

It was not possible to compute the eigenvalues (λ, μ, η) with the minimal values of $|\eta|$ by Algorithms 1 and 3. The difficulty is that these eigenvalues are highly interior. If, instead, we aim for the eigenvalues closest to $(0, 0, 0)$, then the Jacobi–Davidson method performs well with the parameter values indicated above but by solving the correction equations approximately, particularly by employing 10 steps of the generalized minimal residual method with A_i^{-1} as the preconditioner for the ith equation for $i = 1, 2, 3$. The method converges to 50 eigenvalues after 119 subspace updates in 82 s. All but three of the 50 eigenvalues closest to $(0, 0, 0)$ are among the converged eigenvalues, and the remaining eigenvalues converged after a few more iterations. This shows that the Jacobi–Davidson method is capable of locating the eigenvalues closest to a prescribed point, even if these eigenvalues are interior ones.

6 | CONCLUDING REMARKS

We have introduced a Jacobi–Davidson method (Algorithm 1) and a subspace iteration method (Algorithm 3) that restarts the subspace at every iteration for the three-parameter eigenvalue problem. MATLAB implementations are available in package `MultiParEig`.[19] The Jacobi–Davidson method is especially well suited to locate eigenvalues close to a prescribed target. This method seems to perform well in practice also to locate eigenvalues (λ, μ, η) whose η components are close to a prescribed target, whereas the proposed subspace iteration method is specifically designed for this task. Numerical experiments indicate that when the eigenvalues are targeted based on their η components, both methods are very good at locating exterior eigenvalues and mildly interior eigenvalues, but both methods struggle to compute interior eigenvalues.

Based on the numerical experiments, it is not possible to draw a clear conclusion regarding the efficiency of the methods in comparison to each other. In some of the numerical results reported, the Jacobi–Davidson method exhibits better performance in terms of efficiency; in others, the subspace iteration method appears better. To this end, the choice of the parameters, such as the thresholds for the residuals of the Ritz pairs and maximal subspace dimensions, plays an important role.

ACKNOWLEDGEMENTS

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APPENDIX

PROOF OF THEOREM 1

We will only consider the configuration $(\rho, \sigma) = (0, 0)$, as the other three configurations can be treated similarly. Inspired by the work of Cohl and Volkmer,[26] we introduce

$$g(z) := |(\xi - b)(\xi - c)|^{1/2}.$$

We can now write (7) as a three-parameter Sturm–Liouville eigenvalue problem

$$\left(g(\xi_1)X_1'\right)' + \frac{1}{g(\xi_1)} \left(\lambda + \mu\xi_1 + \eta\xi_1^2\right) X_1 = 0, \quad \gamma < \xi_1 < c,$$

$$\left(g(\xi_2)X_2'\right)' + \frac{1}{g(\xi_2)} \left(\lambda + \mu\xi_2 + \eta\xi_2^2\right) X_2 = 0, \quad c < \xi_2 < b,$$

$$\left(g(\xi_3)X_3'\right)' + \frac{1}{g(\xi_3)} \left(\lambda + \mu\xi_3 + \eta\xi_3^2\right) X_3 = 0, \quad b < \xi_3 < \beta.$$

Next, we introduce the elliptic integral

$$G(s) := \int_\gamma^s \frac{d\sigma}{g(\sigma)},$$

which is an increasing absolutely continuous function, and apply the substitution $t_i = G(\xi_i), u_i(t_i) = X_i(\xi_i)$ for $i = 1, 2, 3$. This gives rise to

$$u_1'' + (\lambda + \mu\phi(t_1) + \eta\phi(t_1)^2)u_1 = 0, \quad T_0 < t_1 < T_1,$$
$$u_2'' - (\lambda + \mu\phi(t_2) + \eta\phi(t_2)^2)u_2 = 0, \quad T_1 < t_2 < T_2, \quad \text{(A1)}$$
$$u_3'' + (\lambda + \mu\phi(t_3) + \eta\phi(t_3)^2)u_3 = 0, \quad T_2 < t_3 < T_3,$$

where $T_0 = G(\gamma) = 0, T_1 = G(c), T_2 = G(b), T_3 = G(\beta)$, and $\phi : [T_0, T_3] \to [\gamma, \beta]$ is the inverse function of G. It can be shown that (A1) is a right definite problem due to theorem 3.6.2 in the work of Volkmer.[27] Specifically, let us consider the corresponding determinant function (see, e.g., the work of Atkinson and Mingarelli[4]) given by

$$\delta_0(t_1, t_2, t_3) = \begin{vmatrix} 1 & \phi(t_1) & \phi(t_1)^2 \\ -1 & -\phi(t_2) & -\phi(t_2)^2 \\ 1 & \phi(t_3) & \phi(t_3)^2 \end{vmatrix}$$

$$= (-1)(\phi(t_2) - \phi(t_1))(\phi(t_3) - \phi(t_1))(\phi(t_3) - \phi(t_2)),$$

where $T_0 \le t_1 \le T_1 \le t_2 \le T_2 \le t_3 \le T_3$. One can verify that $\delta_0(t_1, t_2, t_3) < 0$ for all $t_1 < t_2 < t_3$. Since δ_0 is of constant sign on a dense subset of $[T_0, T_1] \times [T_1, T_2] \times [T_2, T_3]$, it follows from theorem 3.6.2 in the work of Volkmer[27] that the problem is right definite. Hence, theorems 3.5.1 and 3.5.2 in the work of Volkmer[27] imply that the Klein oscillation theory holds for the problem. In particular, all eigenvalues are real, and for each triple of nonnegative integers (j_1, j_2, j_3), there exists exactly one eigenvalue (λ, μ, η) such that the corresponding eigenfunction $u_i(t_i)$ has exactly j_i zeros on (T_{i-1}, T_i) for $i = 1, 2, 3$.